

Def 3.1.1 (monoidal cate) :

A monoidal cate is a quintuple  $(\mathcal{C}, \otimes, a, I, \nu)$  where  $\mathcal{C}$  is a cate,  $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$  is a bifunctor  
called the tensor product bifunctor (def of bifunctor, see rothman ex 2.35),  $a: (- \otimes -) \otimes - \xrightarrow{\sim} - \otimes (- \otimes -)$  is  
a natural isom.

$$(2.1) \quad a_{x,y,z}: (x \otimes y) \otimes z \xrightarrow{\sim} x \otimes (y \otimes z), \quad x, y, z \in \mathcal{C}.$$

called the associativity constraint,  $I \in \mathcal{C}$ , an obj of  $\mathcal{C}$ , and  $\nu: I \otimes I \xrightarrow{\sim} I$  is an isom, and satisfy  
the following two axioms.

1. Pentagon axiom The diagram:

$$(2.2) \quad \begin{array}{ccc} & ((w \otimes x) \otimes y) \otimes z & \\ & \swarrow a_{w,x,y} \otimes z & \searrow a_{w,x,y \otimes z} \\ (w \otimes (x \otimes y)) \otimes z & & (w \otimes x) \otimes (y \otimes z) \\ \downarrow a_{w,x \otimes y,z} & & \downarrow a_{w,x,y \otimes z} \\ w \otimes ((x \otimes y) \otimes z) & \xrightarrow{id_w \otimes a_{x,y,z}} & w \otimes (x \otimes (y \otimes z)) \end{array}$$

2. The unit axiom: The functors:

$$(2.3) \quad L_1: X \mapsto 1 \otimes X \quad \text{and} \quad \text{are autoequivalence of } \mathcal{C}.$$

$$(2.4) \quad R_1: X \mapsto X \otimes 1.$$

We say the pair  $(I, \nu)$  The unit obj of  $\mathcal{C}$ .

Why call it monoidal cate?

Reason: The set of isom classes of objs is a monoidal cate indeed has a natural structure of monoid, with  
multiplication  $\otimes$  and unit  $I$ .

2.1.4: A monoidal subcate of a monoidal cate  $(\mathcal{C}, \otimes, a, I, \nu)$  is a quintuple  $(D, \otimes, a, I, \nu)$ , where  
 $D \subset \mathcal{C}$  is a subcate closed under the tensor product of objs and morphisms and containing  $I, \nu$ .

2.1.5:  $\mathcal{C}^{\text{op}} := (\mathcal{C}^{\text{op}}, \otimes^{\text{op}}, 1, a^{\text{op}}, \nu)$ . As a category,  $\mathcal{C}^{\text{op}} = \mathcal{C}$ . Its tensor product:  $x \otimes^{\text{op}} y := y \otimes x$   
and associativity constraint of  $\mathcal{C}^{\text{op}}$  is  $a_{x,y,z}^{\text{op}} = a_{z,y,x}^{-1}: z \otimes (y \otimes x) \xrightarrow{\sim} (z \otimes y) \otimes x$ .  $\mathcal{C}^{\text{op}}$  is a

monoidal cate:

pentagon axiom:

$$\begin{array}{ccc}
 & \text{By (2.2)} & \\
 Z \otimes (Y \otimes (X \otimes W)) & LHS = (\text{id}_W \otimes a_{X,Y,Z}^{\text{op}}) (a_{W,X \otimes Y,Z}^{\text{op}}) (a_{W,X,Y}^{\text{op}} \otimes \text{id}_Z) & \\
 \downarrow \text{id}_Z \otimes a_{Y,X,W}^{-1} & RHS = (a_{W,Y,X}^{\text{op}} \otimes \text{id}_Z) \circ (a_{W \otimes X,Y,Z}^{\text{op}}) & \\
 Z \otimes ((Y \otimes X) \otimes W) & (Z \otimes Y) \otimes (X \otimes W) & \\
 \downarrow a_{Z,Y \otimes X,W}^{-1} & \downarrow a_{Z,Y,X \otimes W}^{-1} & \\
 ((Z \otimes (Y \otimes X)) \otimes W) & \xrightarrow{a_{Z,Y,X \otimes W}^{-1} \otimes \text{id}_W} ((Z \otimes Y) \otimes X) \otimes W &
 \end{array}$$

a.1.b Caution!  $\mathcal{C}^v$  (dual cate of  $\mathcal{C}$ )  $\neq \mathcal{C}^{\text{op}}$  (opposite cate of  $\mathcal{C}$ )

obtained from  $\mathcal{C}$  by reversing arrows.  $\mathcal{C}^v$  is a monoidal cate of  $\mathcal{C}$ .

$L_1(l_x)$ :

$(\mathcal{C}, \otimes, a, 1, \iota)$ : monoidal cate. Define natural isom:

$$(2.5) \quad l_x : 1 \otimes X \longrightarrow X \quad \text{and} \quad r_x : X \otimes 1 \longrightarrow X$$

in such a way s.t.  $L_1(l_x)$  and  $R_1(r_x)$  are equal, resp. to the composition.

$$(2.6) : L_1(l_x) \leftarrow \text{id}_1 \otimes l_x = 1 \otimes (l \otimes X) \xrightarrow{a_{1,1,X}^{-1}} (1 \otimes 1) \otimes X \xrightarrow{\iota \otimes \text{id}_X} 1 \otimes X$$

$$(2.7) : R_1(r_x) \leftarrow r_x \otimes \text{id}_1 = X \otimes 1 \otimes 1 \xrightarrow{a_{X,1,1}^{-1}} X \otimes (1 \otimes 1) \xrightarrow{\text{id}_X \otimes \iota} X \otimes 1.$$

$$\text{i.e. } \text{id}_1 \otimes l_x = (1 \otimes \text{id}_X) \circ l_{1,X}$$

$$r_x \otimes \text{id}_1 = (\text{id}_X \otimes 1) \circ l_{X,1,1}.$$

2.2.1: Isom (2.5) are called left and right unit constraints.

2.2.2: For any obj  $X \in \mathcal{C}$ ,  $\exists$  equalities:

$$(2.8) \quad l_{1 \otimes X} = \text{id}_1 \otimes l_X \quad \text{and} \quad r_{X \otimes 1} = r_X \otimes \text{id}_1.$$

Pf: If follows from the naturality of unit constraints:

$$\begin{array}{ccc}
 1 \otimes (1 \otimes X) & \xrightarrow{\text{id}_1 \otimes l_X} & 1 \otimes X \\
 l_{1 \otimes X} \downarrow & & \downarrow l_X \\
 1 \otimes X & \xrightarrow{l_X \text{ is isom}} & X
 \end{array}
 \xrightarrow{\text{nat}} \text{id}_1 \otimes l_X = l_{1 \otimes X}$$

Similarly

$$\begin{array}{ccc}
 (X \otimes 1) \otimes 1 & \xrightarrow{r_X \otimes \text{id}_1} & X \otimes 1 \\
 r_{X \otimes 1} \downarrow & & \downarrow r_X \\
 X \otimes 1 & \xrightarrow{r_X} & X
 \end{array}$$

2.2.3 The "triangle" diagram:

$$(X \otimes I) \otimes Y \xrightarrow{\alpha_{X,I,Y}} X \otimes (I \otimes Y)$$

$$\begin{array}{ccc} & & \\ r_X \otimes id_Y & \searrow & id_X \otimes l_Y \\ & X \otimes Y & \swarrow \end{array}$$

Pf: Consider

$$\begin{array}{ccccc} ((X \otimes I) \otimes I) \otimes Y & \xrightarrow{\alpha_{X,I,I} \otimes id_Y} & (X \otimes (I \otimes I)) \otimes Y & & \\ \downarrow \alpha_{X \otimes I, I, Y} & \searrow r_{X \otimes I, I, Y} & \curvearrowleft id_{X \otimes I} \otimes id_Y & & \downarrow \alpha_{X, I \otimes I, Y} \\ & (X \otimes I) \otimes Y & \downarrow \alpha_{X, I, Y} & & \\ & \downarrow r_X \otimes id_{I \otimes Y} & & id_X \otimes (l_I \otimes id_Y) & \\ (X \otimes I) \otimes (I \otimes Y) & \xrightarrow{id_X \otimes (l_I \otimes id_Y)} & X \otimes (I \otimes Y) & \xrightarrow{id_X \otimes id_{(I \otimes Y)}} & X \otimes ((I \otimes I) \otimes Y) \\ \downarrow \alpha_{X, I, I \otimes Y} & \curvearrowleft id_X \otimes l_{I \otimes Y} & & \curvearrowleft id_X \otimes id_{(I \otimes I) \otimes Y} & \\ & X \otimes (I \otimes (I \otimes Y)) & & & \end{array}$$

It suffices to show the commutativity of " $\curvearrowleft$ " (Since  $\forall Y \in \mathcal{C}, Y \cong I \otimes Y$  to  $I \otimes Y$ ).

1° The outside pentagon is commutative.

2° " $\curvearrowleft$ " commutes by naturality of also constraint.

3° " $\curvearrowleft$ " commutes by def of  $r$ .

4° " $\curvearrowleft$ " commutes by 2.2. d. Since  $id_L \otimes l_Y = (l \otimes id_Y) \alpha_{L,I,Y}$   
 $\Downarrow 2.2.2$   
 $l_{I \otimes Y}$ .

$$2.2.4: (I \otimes X) \otimes Y \xrightarrow{\alpha_{I,X,Y}} I \otimes (X \otimes Y)$$

$$\begin{array}{ccc} & & \\ (2.12) & \searrow l_X \otimes id_Y & id_X \otimes r_Y \\ & X \otimes Y & \swarrow \end{array}$$

$$(X \otimes Y) \otimes I \xrightarrow{\alpha_{X,Y,I}} X \otimes (Y \otimes I)$$

$$\begin{array}{ccc} & & \\ (2.13) & \searrow r_{X \otimes Y} & id_X \otimes r_Y \\ & X \otimes Y & \swarrow \end{array}$$

Pf: We only show: (2.12). Consider:

$$\begin{array}{ccc}
(x \otimes I \otimes Y) \otimes Z & \xrightarrow{\alpha_{x,I,Y} \otimes \text{id}_Z} & (x \otimes (I \otimes Y)) \otimes Z \\
\downarrow & \swarrow (r_x \otimes \text{id}_Y) \otimes \text{id}_Z & \downarrow \text{id}_{(x \otimes Y)} \otimes \text{id}_Z \\
& (x \otimes Y) \otimes Z & \\
\downarrow \alpha_{x,Y,Z} & & \downarrow \text{id}_{(x \otimes Y)} \otimes \text{id}_Z \\
(x \otimes I) \otimes (Y \otimes Z) & \xrightarrow{r_x \otimes \text{id}_{Y \otimes Z}} & x \otimes (Y \otimes Z) \\
\downarrow & \swarrow \text{id}_{(x \otimes Y \otimes Z)} & \downarrow \text{id}_{(x \otimes (Y \otimes Z))} \\
x \otimes (I \otimes (Y \otimes Z)) & & 
\end{array}$$

↗      ↗      ↗  
↗      ↗      ↗

We need to show the "↗" commutes.

1° The pentagon commutes

2° "↗" naturality of  $\alpha_{\otimes}$  constraint

3° "↗" follows from prop 2.2.3

4° Set  $x = I$  and applying  $L_i^{-1}$  to the lower right triangle

Cor 2.2.5: In any monoidal cate  $l_1 = r_1 = \nu$ .

Pf: Set  $y = x = I$  in (2.2.1). Then  $l_1 \otimes \text{id}_1 = l_1 \otimes l_1 \alpha_{I,I,I} \stackrel{(2.2.8)}{=} (\text{id}_1 \otimes l_1) \circ \alpha_{I,I,I}$ .

Set  $x = y = I$  in (2.2.10) Then  $r_1 \otimes \text{id}_1 = (\text{id}_1 \otimes l_1) \alpha_{I,I,I}$

By def of unit constraint  $(\text{id}_1 \otimes l_1) \circ \alpha_{I,I,I} = l_1 \otimes \text{id}_1 \Rightarrow r_1 \otimes \text{id}_1 = l_1 \otimes \text{id}_1 = \nu \otimes \text{id}_1$   
 Since  $R$  is an autoequivalence. (apply  $R_i^{-1}$ )  
 $\Rightarrow r_1 = l_1 = \nu$ .

2.2.6: The unit obj in a monoidal cate is unique up to a unique bim.

Pf: Let  $(I, \nu)$  and  $(I', \nu')$  be two unit objs. Let  $(r, l)$  and  $(r', l')$  be the corresponding unit constraints. Then we have the norm:  $\eta := l'_1 \circ (r'_1)^{-1} : I \xrightarrow{\sim} I'$

Claim:  $\eta$  maps  $\nu$  to  $\nu'$

We want to show:

$$\begin{array}{ccc}
 1 \otimes 1 & \xrightarrow{\quad r = r_1 \quad} & \text{The whole diag com} \Rightarrow \eta \text{ maps } \iota \text{ to } \iota' \\
 \downarrow \text{id}_1 \otimes \eta & & \text{The rectangle com since } r' \text{ is natural born} \\
 1 \otimes 1' & \xrightarrow{r'_1} & 1 \\
 \downarrow \eta \otimes \text{id}_{1'} & & \downarrow \iota \\
 1' \otimes 1' & \xrightarrow{r'_1} & 1' \\
 \end{array}$$

We want to show:

$$\begin{array}{ccccc}
 x \otimes (1 \otimes 1') & \xrightarrow{\quad \text{id}_x \otimes r'_1 \quad} & & & \Rightarrow r'_1 \circ (\text{id}_1 \otimes l_{1'}) \circ (\text{id}_x \otimes r'_1)^{-1} \\
 & \swarrow \text{id}_x \otimes l_{1'} & \xrightarrow{r' \times 1} & & = r_1 \\
 & & x \otimes 1' & & \\
 & \downarrow r_x \otimes r_{1'} & \xrightarrow{r_x} & & \downarrow r_x \\
 & x \otimes 1' & \xrightarrow{r'_x} & & x
 \end{array}$$

" ↗ " naturality of  $r'$

" ↗ " prop 2.2.3. + prop 2.2.4

It remains to show  $\eta$  is the only born with the property.

It suffices to show if  $b: 1 \rightsquigarrow 1$  the born at the diag commutes.

$$\begin{array}{ccc}
 (2.15) \quad 1 \otimes 1 & \xrightarrow{b \otimes b} & 1 \otimes 1 \\
 \downarrow & & \downarrow \\
 1 & \xrightarrow{b} & 1
 \end{array}$$

Then  $b = \text{id}_1$ . To see this, it suffices to note: for any morphism  $c: 1 \rightarrow 1$  the diag commutes

$$\begin{array}{ccc}
 (2.16) \quad 1 \otimes 1 & \xrightarrow{c \otimes \text{id}_1} & 1 \otimes 1 \\
 \downarrow r_1 = \iota & & \downarrow r = r_1 \\
 1 & \xrightarrow{c} & 1
 \end{array}$$

(since  $\iota = r_1$ ,  $r_1$  is natural born.)

$$b \circ \iota = b \circ (b \otimes b) = b \circ (b \otimes \text{id}_1) \Rightarrow b \circ b = b \otimes \text{id}_1 \Rightarrow b = \text{id}_1.$$

2.2.8: (a traditional def) A monoidal cate is a sextuple  $(\mathcal{C}, \otimes, a, \iota, r)$  satisfying the pentagon axiom

2.2.9 and triangle axiom (2.10)

2.2.9: ① prop 2.2.6 implies for a triple  $(\mathcal{C}, \otimes, a)$  satisfying a pentagon axiom, being a monoidal cate is a property and not a structure

② Furthermore, one can show the commutativity of the triangle implies implies that in a monoidal cate,

one can identify  $1 \otimes X$  and  $X \otimes 1$  with  $X$  using the unit constraints, and assume that the unit morphism are identity.

Take  $\gamma = 1$  in (2.10).

$$(X \otimes 1) \otimes 1 \longrightarrow X \otimes (1 \otimes 1) \Rightarrow r_{X \otimes 1} = R_1(r_X)$$

$\swarrow \quad \searrow$

$$X \otimes 1 \qquad \qquad R_1 \text{ is autoequivalence} \Rightarrow \text{we can identify } X \otimes 1 \text{ with } X.$$

Prop 2.10: Let  $\mathcal{C}$  be a monoidal cate. Then  $\text{End}_{\mathcal{C}}(1)$  is a comonoid under composition. Further more.

$$f \otimes g = v^{-1} \circ (f \circ g) \circ v \text{ for all } f, g \in \text{End}_{\mathcal{C}}(1).$$

Pf: By naturality of unit constraint:

$$f \otimes \text{id}_1 = v^{-1} \circ f \circ r_1. \quad \text{id}_1 \otimes g = v^{-1} \circ g \circ v.$$

$$\text{Since } r_1 = v_1 = v. \Rightarrow$$

$$f \otimes g = (f \otimes \text{id}_1) \circ (\text{id}_1 \otimes g) = v^{-1} \circ (f \circ g) \circ v \Rightarrow f \otimes g = g \otimes f \text{ unit: } v,$$

$$g \otimes f = (\text{id}_1 \otimes f) \circ (g \otimes \text{id}_1) = v^{-1} \circ (f \circ g) \circ v \Rightarrow \text{End}_{\mathcal{C}}(1) \text{ is comonoid}$$

$$f \otimes g = v^{-1} \circ (g \circ f) \circ v \Rightarrow f \circ g = g \circ f$$

2.3: Examples:

2.3.1: The cate Sets form a monoidal cate.

⊗ tensor product: Cartesian product;

② unit obj: one element set;

Caution: The Cartesian product isn't associative, i.e. Sets is not strict.

$$\oplus a: (A \times B) \times C \longrightarrow A \times (B \times C)$$

$$((a, b), c) \mapsto (a, (b, c))$$

$$\oplus l: 1 \times A \longrightarrow A. \quad r: \dots \quad \begin{array}{l} \text{pentagon axiom} \\ \text{triangle axiom} \end{array}$$

$$(1, a) \mapsto a.$$

2.3.2: Any additive cate. is monoidal,

⊗ tensor product: direct sum  $\oplus$

② unit obj: zero obj

$$\oplus a: (A \oplus B) \oplus C \cong A \oplus (B \oplus C)$$

$$(a, b), c \mapsto (a, (b, c))$$

$$\oplus l: 1 \oplus A \longrightarrow A \quad (1, a) \mapsto a.$$

2.3.3: The cate  $\text{rk-Vec}$  of all  $\text{rk-}\mathbb{V}\text{s. where}$

- ①  $\otimes = \otimes_{\mathbb{R}}$ ,
- ②  $1 = \mathbb{R}$ .
- ③  $a, l, r$  are obvious.

More generally, if  $R$  is a com unital ring, then  $R\text{-Mod}$  is a monoidal cate. (not strict)

2.3.4: Let  $G$  be a gp. The cate of  $\text{Rep}_{\mathbb{R}}(G)$  of all rep of  $G$  over  $\mathbb{R}$  is a monoidal cate., with  $\otimes$  being the tensor product of representations,

1. For a rep  $V$ , we denote  $\rho_V : G \rightarrow GL(V)$ .

Then  $V \otimes W$  corresponds to  $\rho_{V \otimes W}(g) := \rho_V(g) \otimes \rho_W(g)$ ,

2. The unit obj: Trivial rep  $1 = \mathbb{R}$ .

2.3.5: Let  $G$  be an affine (pro)alg group over  $\mathbb{R}$ . (an affine algebraic variety with a group structure, s.t the mult map and inverse map are regular (also called a morphism)) is a function between the varieties that's given locally by polynomials. The cate of  $\text{Rep}(G)$ , of algebraic representation of  $G$  over  $\mathbb{R}$  are monoidal categories  
the mult and counit are given similarly

The  $\text{Rep}_{\mathbb{C}_J}$  of a Lie dg  $\mathcal{G}$ . The tensor product:  $\rho_{V \otimes W}(a) = \rho_V(a) \otimes id_W + id_V \otimes \rho_W(a)$   
( $\rho_V : \mathcal{G} \rightarrow gl(V)$ )

The unit: 1-dim rep of  $\mathcal{G}$ , with zero action of  $\mathcal{G}$  we  $\rho_1(a) = 0$

2.3.6: Let  $G$  be a monoid.  $A$  be an abelian group. Let  $\mathcal{E}_G = \mathcal{E}_G(A)$  be the cate where obj's  $S_g$  are labelled by elements of  $G$  (so there's one element in each isom class). Define  $\text{Hom}_{\mathcal{E}_G}(S_{g_1}, S_{g_2}) = \emptyset$  if  $g_1 \neq g_2$ .

and  $\text{Hom}_{\mathcal{E}_G}(S_g, S_g) = A$ .

- ① tensor product of obj  $S_g \otimes S_h = S_{gh}$       tensor product of morphism:  $a \otimes b = ab$
- ② unit obj  $S_1$ .

$\Rightarrow \mathcal{E}_G$  being a monoidal cate with the ~~isomorphism~~ being the identity

Note: in  $\mathcal{E}_G(A)$ ,  $X \otimes Y \neq Y \otimes X$  (for example  $S_{gh}$  may not be equal to  $S_{hg}$ )

A "mean" version:  $\text{rk-}\mathbb{V}\text{ec}_G$ :  $G$ -graded  $\mathbb{V}\text{s. over } \mathbb{R}$ .

Vector space  $V \in \text{Vec}_G$  has a decomposition  $V = \bigoplus_{g \in G} V_g$ . Morphism in  $\text{Vec}_G$  preserves the grading

1° tensor product :  $V \otimes W$  as vector space tensor product,

$$\text{and } (V \otimes W)_g = \bigoplus_{\substack{x,y \in G \\ xy=g}} V_x \otimes W_y$$

2° unit obj  $1$  by  $1_1 = \mathbb{K}$  and  $1_g = 0$  for  $g \neq 1$ .

3° a.,  $(U \otimes V) \otimes W \xrightarrow{\alpha_{U,V,W}} U \otimes (V \otimes W)$   $\Rightarrow \text{Vec}_G$  is a monoidal cate

$$\iota : 1 \otimes 1 \rightarrow 1.$$

In  $\text{Vec}_G$ , we have parawire obj  $S_g$  defined by  $(S_g)_x = \mathbb{K}$  if  $x=g$  and  $(S_g)_x = 0$  otherwise, by def.

$S_g \otimes S_h = S_{gh}$ . Thus :  $\mathcal{C}_G(\mathbb{K})$  is a non-full monoidal subcate of  $\text{Vec}_G$   
 $(\text{Zero morphism miss})$

This subcate can be viewed as a "linear basis" of  $\text{Vec}_G$ . As any obj of  $\mathbb{K}\text{-Vec}_G$  is iso to a direct sum of  
 objs  $S_g$  with non-negative integer multiplicities