# ELEMENTARY DIFFERENCES AMONG FINITE LEVELS OF THE ERSHOV HIERARCHY

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ABSTRACT. We study the differences among finite levels of the Ershov hierarchies. We also give a brief survey on the current state of this area. Some questions are raised.

### 1. Preliminary

Putnam [9] is the first one who introduced the n-r.e. sets.

**Definition 1.1.** (i) A set A is n-r.e. if there is a recursive function  $f: \omega \times \omega \rightarrow \omega$  so that for each m,

-f(0,m) = 0.-  $A(m) = \lim_{s} f(s,m).$ 

 $-|\{s|f(s+1,m) \neq f(s,m)\}| \le n.$ 

• A Turing degree is n-r.e. if it contains an n-r.e. set.

We use  $D_n$  to denote the collection of *n*-r.e. degrees. For simplicity, we redefine  $D_0 = D_1$  which is a little unusual.

For other recursion notations, please refer to Soare [13].

In this paper, we work in the partially ordered language,  $\mathcal{L}(\leq)$ , through out.  $\mathcal{L}(\leq)$  includes variables a, b, c, x, y, z, ... and a binary relation  $\leq$  intended to denote a partial order. Atomic formulas are  $x = y, x \leq y$ .  $\Sigma_0$  formulas are built by the following induction definition.

- Each atomic formula is  $\Sigma_0$ .
- $\neg \psi$  for some  $\Sigma_0$  formula  $\psi$ .
- $\psi_1 \lor \psi_2$  for two  $\Sigma_0$  formula  $\psi_1, \psi_2$ .
- $\psi_1 \wedge \psi_2$  for two  $\Sigma_0$  formula  $\psi_1, \psi_2$
- $\psi_1 \implies \psi_2$  for two  $\Sigma_0$  formula  $\psi_1, \psi_2$ .

A formula  $\varphi$  is  $\Sigma_1$  if it is of the form  $\exists x_1 \exists x_2 ... \exists x_n \psi(x_1, x_2, ..., x_n)$  for some  $\Sigma_0$  formula  $\psi$ .

For each  $n \geq 1$ , a formula  $\varphi$  is  $\Pi_n$  if it is the form  $\neg \psi$  for some  $\Sigma_n$  formula  $\psi$  and a formula  $\varphi$  is  $\Sigma_{n+1}$  if it is the form  $\exists x_1 \exists x_2 ... \exists x_m \psi(x_1, x_2, ..., x_m)$  for some  $\Pi_n$  formula  $\psi$ .

A sentence is a formula without free variables.

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Given two structures  $\mathfrak{A}(A, \leq_A)$  and  $\mathfrak{B}(B, \leq_B)$  for  $\mathcal{L}(\leq)$ , we say that  $\mathfrak{A}(A, \leq_A)$ is a substructure of  $\mathfrak{B}(B, \leq_B)$ , write  $\mathfrak{A}(A, \leq_A) \subseteq \mathfrak{B}(B, \leq_B)$ , if  $A \subseteq B$  and the interpretation  $\leq_A$  is a restriction to A of  $\leq_B$ .

**Definition 1.2.** For  $n \ge 0$ . Given structures  $\mathfrak{A}(A, \leq_A)$  and  $\mathfrak{B}(B, \leq_B)$  for  $\mathcal{L}(\leq)$ .

(i) We say that  $\mathfrak{A}(A, \leq_A)$  is a  $\Sigma_n$  substructure of  $\mathfrak{B}(B, \leq_B)$ , write  $\mathfrak{A}(A, \leq_A) \preceq_{\Sigma_n} \mathfrak{B}(B, \leq_B)$ , if  $\mathfrak{A}(A, \leq_A) \subseteq \mathfrak{B}(B, \leq_B)$  and for all  $\Sigma_n$  formulas  $\varphi(x_1, x_2, ..., x_n)$ and any  $a_1, a_2, ..., a_n \in A$ ,

 $\mathfrak{A}(A, \leq_A) \models \varphi(a_1, a_2, ..., a_n)$  if and only if  $\mathfrak{B}(B, \leq_B) \models \varphi(a_1, a_2, ..., a_n)$ .

(ii) We say that  $\mathfrak{A}(A, \leq_A)$  is  $\Sigma_n$ -elementary-equivalent to  $\mathfrak{B}(B, \leq_B)$ , write  $\mathfrak{A}(A, \leq_A) \equiv_{\Sigma_n} \mathfrak{B}(B, \leq_B)$ , if for all  $\Sigma_n$  sentences  $\varphi$ ,

 $\mathfrak{A}(A, \leq_A) \models \varphi$  if and only if  $\mathfrak{B}(B, \leq_B) \models \varphi$ .

In this paper, we study the model theoretical properties of  $\Delta_2^0$  Turing degrees as the structure  $\mathcal{D}(\leq \mathbf{0}') = (\mathcal{D}(\leq \mathbf{0}'), \leq)$  of  $\mathcal{L}(\leq)$ . We are interested in various substructure of  $\mathcal{D}(\leq \mathbf{0}')$ , particularly, the structures of *n*-r.e. degrees  $\mathcal{D}_n = (\mathcal{D}_n, \leq)$ .<sup>1</sup> For two degrees **a** and **b** in  $\mathcal{D}_n$  (or  $\mathcal{D}(\leq \mathbf{0}')$ ), we use  $\mathbf{a} \cup \mathbf{b}$  and  $\mathbf{a} \cap \mathbf{b}$  to denote their least upper bound and the largest lower bound (if exists) in  $\mathcal{D}_n$  (or  $\mathcal{D}(\leq \mathbf{0}')$ ) respectively.

For more model theoretic facts, please refer to [7].

2. Elementary difference among Ershov hierarchies

Comparing the structure difference between Ershov hierarchies has a long history beginning with Cooper(1970's) and Lachlan's (1968) unpublished work. They proved the following theorem.

**Theorem 2.1** (Lachlan, Cooper). (i) For each  $n \ge 1$ ,  $D_n \subset D_{n+1}$ .

(ii) For each non-recursive n + 1-r.e. degree  $\mathbf{a}$ , there is a non-recursive n-r.e. degree  $\mathbf{b} \leq \mathbf{a}$ .

For any  $\Sigma_1$ -sentence  $\varphi$ ,  $\mathcal{D}_n$  or  $\mathcal{D}(\leq \mathbf{0}')$  satisfies  $\varphi$  if and only if  $\varphi$  is consistent with the theory of partial orderings (see, for example, some exercises in Soare [13]). Therefore,

**Theorem 2.2** (Folklore). For all  $n \in \omega$ ,  $\mathcal{D}_n \equiv_{\Sigma_1} \mathcal{D}(\leq \mathbf{0}')$ .

Thus elementary differences would not occur at the  $\Sigma_1$ -level.

By improving a technique due to Spector, Sacks proved the following result.

**Theorem 2.3** (Sacks [10]). There is a  $\Delta_2^0$  minimal degree.

Comparing with Theorem 2.1, the elementary difference between  $\mathcal{D}_n$  and  $\mathcal{D}(\leq \mathbf{0}')$  shows up at  $\Sigma_2$ -level.

The elementary difference between  $\mathcal{D}_1$  and  $\mathcal{D}_n(n > 1)$  was first revealed at  $\Sigma_3$ -level by Arslanov [2] who showed that for every element **a** in  $\mathcal{D}_n$ , there is an element  $\mathbf{b} \in \mathcal{D}_n$  of which the supreme is  $\mathbf{0}'$ , whereas in  $\mathcal{D}_1$  this is not true due to Cooper and Yates. Later many differences at  $\Sigma_2$ -level were discovered, for example, the following pair of theorems offers perhaps the clearest order-theoretic difference:

 $\mathbf{2}$ 

<sup>&</sup>lt;sup>1</sup>We use "1-r.e." to denote "r.e."

**Theorem 2.4** (Sacks[11]).  $\mathcal{D}_1$  is dense.

**Theorem 2.5** (Cooper, Harrington, Lachlan, Lempp, Soare[5]). For each natural number n > 1, there is a maximal element in  $\mathcal{D}_n$ .

So the following results can be obtained.

**Corollary 2.6.** For each natural number n > 1,  $\mathcal{D}_1 \not\equiv_{\Sigma_2} \mathcal{D}_n$ .

A further question is how difference between  $\mathcal{D}_n$  and  $\mathcal{D}_{n+m}$  for n > 1. Downey formulated the following ambitious question which is now known as Downey Conjecture.

Conjecture 2.7 (Downey [6]). For each n > 1 and  $k \ge 0$ ,  $\mathcal{D}_n \equiv_{\Sigma_k} \mathcal{D}_{n+m}$ .

Though Downey Conjecture looks too optimal to be true, it remained open more than fifteen years. The difficulty of Conjecture 2.7 lies in the technique used in the local theory of  $\mathcal{D}_n$ . Usually one can generalize a (local) result in  $\mathcal{D}_2$  to  $\mathcal{D}_n$  without any difficult.

Recently, Arslanov, Kalimullin and Lempp announced a negative solution to Conjecture 2.7. They proved the following result.

**Theorem 2.8** (Arslanov, Kalimullin, Lempp [3]).  $\mathcal{D}_2 \not\equiv_{\Sigma_2} \mathcal{D}_3$ .

But the question whether  $\mathcal{D}_n \not\equiv_{\Sigma_2} \mathcal{D}_{n+m}$  is true for some very large numbers n, m still remains open.

3.  $\Sigma_1$ -substructures of  $\mathcal{D}(\leq \mathbf{0}')$ 

As we have seen that  $\mathcal{D}_n \equiv_{\Sigma_1} \mathcal{D}(\leq \mathbf{0}')$  (Theorem 2.2), it is natural to ask whether  $\mathcal{D}_n \preceq_{\Sigma_1} \mathcal{D}(\leq \mathbf{0}')$ . This was eventually negatively answered by Slaman in 1983.

**Theorem 3.1** (Slaman). (i) There are r.e. sets A, B and C and a  $\Delta_2^0$  set E such that

 $-\emptyset <_T E \leq_T A;$ 

$$-C \not\leq_T B \oplus E;$$

- For all r.e. set  $W \ (\emptyset <_T W \leq_T A \Rightarrow C \leq_T W \oplus B)$ .

(ii) For each natural number  $n \ge 1$ ,  $\mathcal{D}_n \not\preceq_{\Sigma_1} \mathcal{D}(\le \mathbf{0}')$ .

*Proof.* We just show how to deduce (ii) from (i). Take a  $\Sigma_1$  formula

$$\varphi(x_1, x_2, x_3) \equiv \exists e \exists y \exists z (e \le x_1 \land e \ge y \land e \ne y \land z \ge x_2 \land z \ge e \land z \ge x_3).$$

Take the r.e. degrees  $\mathbf{a}, \mathbf{b}, \mathbf{c}$  and a  $\Delta_2^0$  degree  $\mathbf{e}$  as in (i). Fix  $Z = B \oplus E$ .

Then  $\mathcal{D}(\mathbf{0}') \models \varphi(\mathbf{a}, \mathbf{b}, \mathbf{c})$  since  $\mathbf{e} \leq \mathbf{a} \land \mathbf{e} > \mathbf{0} \land \mathbf{z} \not\geq \mathbf{c}$ .

Then for each  $n \ge 1$ ,  $\mathcal{D}_n \not\models \varphi(\mathbf{a}, \mathbf{b}, \mathbf{c})$ . If not, then there is an *n*-r.e. degree  $\mathbf{f} > \mathbf{0}$  so that  $\mathbf{f} \le \mathbf{a}$  and  $\mathbf{f} \cup \mathbf{b} \not\ge \mathbf{c}$ . But, by Theorem 2.1, there is a non-recursive r.e. degree  $\mathbf{w} \le \mathbf{f}$ . So  $\mathbf{w} \cup \mathbf{b} \not\ge \mathbf{c}$ . This is impossible by (i).

Having proved Theorem 3.1, Slaman raised the following conjecture which remained open more than twenty years.

Conjecture 3.2 (Slaman [5]). For each n > 1,  $\mathcal{D}_1 \preceq_{\Sigma_1} \mathcal{D}_n$ ?

Furthermore, Lempp raised the following conjecture.

Conjecture 3.3 (Lempp). For all n > m,  $\mathcal{D}_m \preceq_{\Sigma_1} \mathcal{D}_n$ ?

To solve conjecture 3.2, one possible argument is to build a finite array just as Slaman did in Theorem 3.1. However, by the Cooper and Lachlan observation that every nonrecursive *n*-r.e. degree bounds a nonrecursive r.e. degree, we cannot hope that any *n*-r.e. degree D plays the role of E as in Slaman Theorem.

We first explain that it is necessary to build a complicated formula to refute Slaman's conjecture.

A formula is called positive if it is built by the following induction definition.

- Each atomic formula is positive.
- $\psi_1 \lor \psi_2$  for two positive formula  $\psi_1, \psi_2$ .
- $\psi_1 \wedge \psi_2$  for two positive formula  $\psi_1, \psi_2$ .

A formula is called  $p-\Sigma_1$  if it is the form  $\exists x_1 \exists x_2 ... \exists x_n \varphi(x_1, x_2, ..., x_n)$  for some positive formula  $\varphi$ .

We say that  $\mathfrak{A}(A, \leq_A)$  is a p- $\Sigma_1$  substructure of  $\mathfrak{B}(B, \leq_B)$ , write  $\mathfrak{A}(A, \leq_A) \preceq_{p$ - $\Sigma_1} \mathfrak{B}(B, \leq_B)$ , if  $\mathfrak{A}(A, \leq_A) \subseteq \mathfrak{B}(B, \leq_B)$  and for all p- $\Sigma_1$  formulas  $\varphi(x_1, x_2, ..., x_n)$  and any  $a_1, a_2, ..., a_n \in A$ ,

 $\mathfrak{A}(A, \leq_A) \models \varphi(a_1, a_2, ..., a_n)$  if and only if  $\mathfrak{B}(B, \leq_B) \models \varphi(a_1, a_2, ..., a_n)$ .

We have the following proposition

**Proposition 3.4.**  $\mathcal{D}_n \leq_{p \to \Sigma_1} \mathcal{D}_m$  for all  $n \leq m$ . Furthermore,  $(D_1, \leq, \cup, \cap) \leq_{p \to \Sigma_1} (D_n, \leq, \cup, \cap)$  for all n > 1.

Thus to refute Slaman Conjecture, it is necessary to consider some negative statement.

Eventually we obtained the following formula.

 $\varphi(x_1, x_2, x_3, x_4) \equiv \exists d \exists g (d \le x_1 \land d \le x_4 \land g \ge x_2 \land g \ge d \land x_3 \le g).$ 

The solution to Conjecture 3.2 follows from the following technical result:

**Theorem 3.5** (Yang and Yu [15]). There are r.e. sets A, B, C and E and a d.r.e. set D such that

- (1)  $D \leq_T A$  and  $D \not\leq_T E$ ;
- (2)  $C \not\leq_T B \oplus D;$
- (3) For all r.e. sets W ( $W \leq_T A \Rightarrow$  either  $C \leq_T W \oplus B$  or  $W \leq_T E$ ).

Assuming Theorem 3.5, we can obtain the following result to refute Slaman conjecture:

**Theorem 3.6** (Yang and Yu [15]). For all n > 1,  $\mathcal{D}_1 \not\preceq_{\Sigma_1} \mathcal{D}_n$ .

*Proof.* Assume n > 1.Let  $\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{d}, \mathbf{e}$  be the degrees of their corresponding sets as in Theorem 3.5. Note all of them except  $\mathbf{d}$  belong to  $\mathcal{D}_1$  and  $\mathbf{d}$  belongs to  $\mathcal{D}_n$ . By Theorem 3.5, just take  $\mathbf{g} = \mathbf{b} \cup \mathbf{d} \in \mathcal{D}_n$ ,

$$\mathcal{D}_n \models \varphi(\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{e}).$$

4

However, by Theorem 3.5 again,

$$\mathcal{D}_1 \models \forall \mathbf{w} \forall \mathbf{g} ((\mathbf{w} \le \mathbf{a} \land \mathbf{g} \ge \mathbf{w} \land \mathbf{g} \ge \mathbf{b} \land \mathbf{w} \not\le \mathbf{e}) \implies \mathbf{c} \le \mathbf{g}).$$

In other words,

$$\mathcal{D}_1 \models \neg \varphi(\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{e}).$$

Although Slaman Conjecture is not true, we can ask where the abnormal parameters refuting the conjecture exist. Inspired by Shore and Slaman [12], we conjecture that each high r.e. degree bounds the four parameters as in Theorem 3.5 so that Slaman conjecture fails. But is there a fragment  $\mathcal{E} \subset \mathcal{D}_1$  so that  $\mathcal{E} \preceq_{\Sigma_1} \mathcal{D}_2$ ? A critical part of the argument used in the proof of Theorem 3.5 is a modification of the construction of Slaman triple. A triple  $(\mathbf{a}, \mathbf{b}, \mathbf{c})$  in  $\mathcal{D}_n$  is called *Slaman triple* if  $\mathbf{0} < \mathbf{a}, \mathbf{c} \leq \mathbf{b}$  and for all non-recursive  $\mathbf{x} \in \mathcal{D}_n$  below  $\mathbf{a}, \mathbf{c} \leq \mathbf{b} \cup \mathbf{x}$ . Shore and Slaman [12] showed that a Slaman-triple can be found below each high r.e. degree in  $\mathcal{D}_1$ . However, Harrington, and Bickford and Mills, showed independently that no low<sub>2</sub> r.e. degree bounds a Slaman triple in  $\mathcal{D}_1$ . Thus it sounds reasonable to conjecture that there is fragment  $\mathcal{E} \subset \mathcal{D}_1$  in which all of elements are low<sub>2</sub> so that  $\mathcal{E} \preceq_{\Sigma_1} \mathcal{D}_2$ . A non-recursive degree  $\mathbf{a} \in \mathcal{D}_n$  is said to be *almost deep* if for each low  $\mathbf{b} \in \mathcal{D}_n, \mathbf{a} \cup \mathbf{b}$  is low. Cholak et al [4] proved that almost deep degrees exist in  $\mathcal{D}_1$ . Hence it is natural to ask whether the almost deep degrees in  $\mathcal{D}_1$  form a  $\Sigma_1$ -substructure of  $\mathcal{D}_2$ .

The last question in this section was raised by Khoussainov.

**Question 3.7** (Khoussainov). For n > 1, is there a function  $f : D_1 \to D_n$  so that for any  $\Sigma_1$ -formula  $\varphi(x_1, ..., x_m)$ ,

$$\mathcal{D}_1 \models \varphi(\mathbf{x_1}, \mathbf{x_2}, ..., \mathbf{x_m}) \text{ iff } \mathcal{D}_n \models \varphi(f(\mathbf{x_1}), f(\mathbf{x_2}), ..., f(\mathbf{x_m})),$$

where  $\mathbf{x_1}, \mathbf{x_2}, ..., \mathbf{x_m}$  range over  $D_1$ ?

#### 4. Definable ideals and filters

Recently, Wang and Yu [14] proved that each non-principal ideal in  $\mathcal{D}_1$  is a  $\Sigma_1$ substructure of  $\mathcal{D}_1$ . But the question whether any non-principal ideal in  $\mathcal{D}_2$  is a  $\Sigma_1$ -substructure of  $\mathcal{D}_2$  is unknown. A set  $\mathcal{A} \subseteq \mathcal{D}_n$  is said to be definable in  $\mathcal{D}_n$  if there is a formula  $\psi$  so that  $\mathbf{a} \in \mathcal{A}$  if and only if  $\mathcal{D}_n \models \psi(\mathbf{a})$ . For  $\mathcal{D}_1$ , by the recent work due to Nies [8], Yang and Yu [16], there are many definable non-principal ideals in  $\mathcal{D}_1$ . A natural question is what about  $\mathcal{D}_2$ ? To construct a non-principal ideal in  $\mathcal{D}_n$ , we just need to take a non-principal ideal  $\mathcal{I}$  in  $\mathcal{D}_1$  and then build a non-principal ideal  $\mathcal{J} = \{\mathbf{b} | \exists \mathbf{a} \in \mathcal{I}(\mathbf{b} \leq \mathbf{a})\}$ . The problem is whether it is definable in  $\mathcal{D}_n$ . We formulate the following questions which we are very interested in.

## Question 4.1. For n > 1, is there a non-trivial definable $\Sigma_1$ -substructure of $\mathcal{D}_n$ ?

From the discussion above, we have seen that the definable ideals play a critical role in the study of global theory. Although there are some non-trivial definable ideals in  $\mathcal{D}_1$ . It is unknown whether there are infinitely many definable ones in  $\mathcal{D}_1$ . For  $\mathcal{D}_2$ , we don't even know whether there is a non-trivial definable ideal in it.

#### YUE YANG AND LIANG YU

Wang also recently studied definable filters in  $\mathcal{D}_1$ . It is unknown whether there is a non-trivial definable filter in  $\mathcal{D}_2$ . We say that a non-zero degree  $\mathbf{a} \in \mathcal{D}_n$  is *cappable* if there is a non-zero degree  $\mathbf{b} \in \mathcal{D}_n$  so that the infimum of them is the recursive degree **0**. Otherwise, **a** is said to be *non-cappable*. A possible candidate of definable filters is the collection of non-cappable degrees in  $\mathcal{D}_2$ . Ambos-Spies et al [1] proved that the collection of non-cappable degrees form a filter in  $\mathcal{D}_1$ . Thus, to construct a definable filter in  $\mathcal{D}_2$ , it suffices to prove that each non-cappable 2-r.e. degree bounds a non-cappable r.e. degree. So the following technique question is raised.

**Question 4.2.** Can each 2-r.e. non-cappable degree compute an r.e. non-cappable degree?

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 $\mathbf{6}$