# A NOTE ON REDUCIBLE CYCLES IN MULTIPARTITE TOURNAMENTS* 

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#### Abstract

T\) is a strong $c$-partite tournament $(c \geq 3)$, then there is a $(k-3)$-reducible $k$-cycle in $T$, for all $k=3,4, \cdots, c$. In this paper we investigate the smallest number of $(k-3)$-reducible $k$-cycles in strong $c$-partite tournaments for $3 \cdot k \cdot c$ and give some related problems.


## 1. Introduction

We assume that the reader is familar with the standard terminology on graphs and digraphs and refer the reader to [2].

A digraph $D=(V(D), A(D))$ is determined by its set of vertices $V(D)$, and its set of arcs $A(D)$. If $x y$ is an arc of a digraph $D$, then we say that $x$ dominates $y$ and write $x \rightarrow y$. More generally, if $A$ and $B$ are two disjoint subdigraphs of $D$ or subsets of $V(D)$ such that every vertex of $A$ dominates every vertex of $B$, then we say that $A$ dominates $B$ and write $A \rightarrow B$. We use $A \Rightarrow B$ to denote the fact that there is no arc leading from $B$ to $A$. By a cycle (path, resp.) we mean a directed cycle (directed path, resp.). A digraph $D$ is strong if for any two vertices $x$ and $y$ there exists a path from $x$ to $y$ and a path from $y$ to $x$ in $D$. A cycle of length $k$ is called a $k$-cycle. A cycle (path, resp.) of a digraph $D$ is Hamiltonian if it includes all the vertices of $D$. A digraph $D$ is pancyclic if it contains a $k$-cycle for all $k$ between 3 and $|V(D)|$. A digraph $D$ is vertex pancyclic if every vertex of $D$ is contained in a $k$-cycle for all $k \in\{3,4, \cdots,|V(D)|\}$. If $S$ is a set of vertices in a digraph $D$, then $D[S]$ is the subgraph induced by $S$.

A c-partite or multipartite tournament is a digraph obtained from a complete $c$-partite graph by substituting each edge with an arc. Let $T$ be a multipartite

[^0]tournament and $v \in V(T)$. We use $V^{c}(v)$ to denote the partite set which $v$ belongs to.

Let $D$ be a digraph and let $k$ be some integer. A cycle $C_{0}$ is $k$-reducible if there are cycles $C_{1}, C_{2}, \cdots, C_{k}$ such that for all $i=0,1, \cdots, k-1$ there is a vertex $w_{i}$ in $C_{i}$ such that $C_{i+1}=C_{i}\left[w_{i}^{+}, w_{i}^{-}\right] w_{i}^{+}$. Let $w \in V(D)$. Then a cycle $C_{0}$ is $(w, k)$-reducible if it is $k$-reducible and $w$ belongs to all the cycles $C_{1}, C_{2}, \cdots, C_{k}$ (i.e., $w_{i} \neq w$ for all $i=0,1, \cdots, k-1$ ).
[1] proves that if $T$ is a strong $c$-partite tournament $(c \geq 3)$, then there is a $k$-cycle in $T$ for all $k=3,4, \cdots, c$. [3] extends this result by showing that if $T$ is a strong $c$-partite tournament $(c \geq 3)$, then there is a $(k-3)$-reducible $k$-cycle in $T$ for all $k=3,4, \cdots, c$. In this paper we investigate the smallest number of $(k-3)$-reducible $k$-cycles in strong $c$-partite tournaments for $3 \cdot k \cdot c$.

Theorem. Let $T$ be a strong c-partite tournament. Then the number of $(k-3)$ reducible $k$-cycles is at least $c-k+1$ for $3 \cdot k \cdot c$. Moreover, the lower bound is best possible.

## 2. Proof of Theorem

Lemma 1 (Yeo [3]). If $T$ is a strong $c$-partite tournament ( $c \geq 3$ ), then there is a $(k-3)$-reducible $k$-cycle in $T$ for all $k=3,4, \cdots, c$.

Lemma 2 (Goddard \& Oellermann [4]). Every vertex of a strong c-partite tournament $(c \geq 3)$ belongs to a cycle which contains vertices from exactly $q$ partite sets for each $q \in\{3,4, \cdots, c\}$.

Lemma 3 (Guo \& Volkmann [5]). Every partite set of a strong c-partite $(c \geq 3)$ tournament has at least one vertex which lies on a $k$-cycle for each $k \in\{3,4, \cdots, c\}$.

The following Lemma 4 is interesting in itself; it generalizes Moon's theorem on vertex pancyclicity in strong tournaments [6].

Lemma 4. Let $T$ be a strong c-partite tournament with partite sets $V_{1}, V_{2}, \cdots$, $V_{c}$. Then for any $V_{i}$ there is a $(k-3)$-reducible $k$-cycle in $T$ which contains at least one vertex of $V_{i}$ for all $k=3,4, \cdots, c$.

Proof. We prove the lemma by induction on $k$. When $k=3$, Lemma 4 holds by Lemma 3. We assume $4 \cdot k \cdot c$ and $V_{i}$ is given. By Lemma 1 , there is a $(k-3)$-reducible $k$-cycle $C_{0}$ in $T$. If $V\left(C_{0}\right) \cap V_{i} \neq \emptyset$, we are done. So we assume that $V\left(C_{0}\right) \cap V_{i}=\emptyset$ and take a $(k-4)$-reducible $(k-1)$-cycle $C_{1}$ in $T$ such that $V\left(C_{1}\right) \cap V_{i}=\emptyset$ (such a cycle exists by the reducibility of cycle $C_{0}$ ). If there is a
vertex $v \in V_{i}$ with $v \nRightarrow V\left(C_{1}\right)$ and $V\left(C_{1}\right) \nRightarrow v$, then since $v$ is adjacent to every vertex in $V\left(C_{1}\right)$ there exists a vertex $u \in V\left(C_{1}\right)$ such that $u^{-} \rightarrow v$ and $v \rightarrow u$. We obtain a $k$-cycle $C_{1}\left[u, u^{-}\right] v u$, which is $(k-3)$-reducible and contains a vertex $v$ of $V_{i}$. Therefore we assume that for each $v \in V_{i}$ either $v \Rightarrow V\left(C_{1}\right)$ or $V\left(C_{1}\right) \Rightarrow v$.

Let $A_{1}=\left\{v \in V_{i} \mid v \Rightarrow V\left(C_{1}\right)\right\}$ and $A_{2}=\left\{v \in V_{i} \mid V\left(C_{1}\right) \Rightarrow v\right\}$. Clearly $A_{1} \cup A_{2}=V_{i}$. Since $V\left(C_{1}\right) \cap V_{i}=\emptyset$, it is easy to see that $A_{1} \rightarrow V\left(C_{1}\right)$ and $V\left(C_{1}\right) \rightarrow A_{2}$. Let $l$ be the length of a shortest path of all $\left(C_{1}, A_{1}\right)$-paths and $\left(A_{2}, C_{1}\right)$-paths. Without loss of generality, we assume that $P=y_{0} y_{1} \cdots y_{l}$ is a $\left(C_{1}, A_{1}\right)$-path of length $l$. If $y_{1} \in V^{c}\left(y_{l}\right)$, then $l \geq 3$ and $y_{1} \in A_{2}$. Let $x \in V\left(C_{1}\right)-$ $V^{c}\left(y_{2}\right)$ be arbitrary. If $x \rightarrow y_{2}$, then the path $x P\left[y_{2}, y_{l}\right]$ is a shorter $\left(C_{1}, A_{1}\right)$-path than $P$, a contradiction. If $y_{2} \rightarrow x$, then the path $y_{1} y_{2} x$ is a shorter $\left(A_{2}, C_{1}\right)$ path than $P$, a contradiction. So $y_{1} \notin V^{c}\left(y_{l}\right)$. Similarly, from the minimality of $l$ we obtain that $V^{c}\left(y_{l}\right) \cap\left\{y_{1}, y_{2}, \cdots, y_{l-1}\right\}=\emptyset$ and $y_{l} \rightarrow\left\{y_{0}, y_{1}, \cdots, y_{l-2}\right\}$. Let $C_{2}=P C_{1}\left[y_{0}^{+l}, y_{0}\right]$. Since $y_{l} \rightarrow V\left(C_{2}\right)-\left\{y_{l-1}\right\}, C_{2}$ is a $(k-3)$-reducible $k$-cycle and contains a vertex $y_{1}$ of $V_{i}$.

This completes the proof of Lemma 4.
Corollary 5 (Moon [6]). Every strong tournament is vertex pancyclic.
Theorem 6. Let $T$ be a strong c-partite tournament. Then the number of $(k-3)$-reducible $k$-cycles is at least $c-k+1$ for $3 \cdot k \cdot c$.

Proof. Let $V_{1}, V_{2}, \cdots, V_{c}$ be the partite sets of $T$ and $3 \cdot k \cdot c$. We prove the theorem by induction on $c$. For $c=k$, the result follows from Lemma 1 .

Suppose now that $c \geq k+1$ and that every strong ( $c-1$ )-partite tournament contains at least $(c-1)-k+1(k-3)$-reducible $k$-cycles. According to Lemma 2, there exists a cycle $C$ that contains vertices from exactly $c-1$ partite sets. $T[V(C)]$ is a strong $(c-1)$-partite tournament, which contains by the induction hypothesis at least $c-k(k-3)$-reducible $k$-cycles. Without loss of generality, let $V_{1}$ be the partite set with $V_{1} \cap V(C)=\emptyset$. By Lemma 4, $T$ contains a $(k-3)$-reducible $k$-cycle $C_{2}^{\prime}$ with $V_{1} \cap V\left(C_{2}^{\prime}\right) \neq \emptyset$. Clearly the $(k-3)$-reducible $k$-cycle $C_{2}^{\prime}$ is different from the $(k-3)$-reducible $k$-cycles in $T[V(C)]$. So $T$ contains at least $c-k+1$ ( $k-3$ )-reducible $k$-cycles.

Corollary 7. Let $T$ be a strong c-partite tournament $(c \geq 3)$. Then there are at least $c-k+1$ pancyclic subgraphs of order $k$ in $T$ for all $k=3,4, \cdots, c$.

Corollary 8 (Goddard \& Oellermann [4]). Let $T$ be a strong c-partite tournament $(c \geq 3)$. Then $T$ contains at least $c-2$ cycles of length 3 .

Corollary 9. Let $T$ be a strong c-partite tournament $(c \geq 3)$. Then $T$ contains at least $\binom{c-1}{2}$ cycles.

Corollary 10 (Moon [6]). Let $T$ be a strong tournament of order $n$. Then $T$ contains at least $n-k+1$ cycles of length $k$ for $3 \cdot k \cdot n$.

Corollary 11 (Moon [6]). Let $T$ be a strong tournament of order $n$. Then $T$ contains at least $\binom{n-1}{2}$ cycles.

The tournament obtained by reversing the arcs of the unique Hamiltonian path in a transitive tournament $T_{n}$ with $n$ vertices is seen to have precisely $n-k+1$ $(k-3)$-reducible $k$-cycles for $3 \cdot k \cdot n$. This example shows that the estimation in Theorem 6 is best possible. We denote the above tournament by $M_{n}$.

We construct a 6 -partite tournament $T_{6}$ with partite sets $V_{i}=\left\{V_{i}\right\}, i=$ $1,2,3,4,5$, and $V_{6}=\left\{v_{6}, v_{7}\right\}$. Let $v_{1} \rightarrow\left\{v_{2}, v_{4}, v_{5}, v_{7}\right\}, v_{2} \rightarrow\left\{v_{3}, v_{4}, v_{5}, v_{6}, v_{7}\right\}$, $v_{3} \rightarrow\left\{v_{1}, v_{4}, v_{5}, v_{6}, v_{7}\right\}, v_{4} \rightarrow\left\{v_{5}, v_{6}\right\}, v_{5} \rightarrow v_{7}, v_{6} \rightarrow\left\{v_{1}, v_{5}\right\}$ and $v_{7} \rightarrow v_{4}$. It is easy to see that $T_{6}$ is strong and $T_{6}$ contains no strong tournament with 6 vertices. This example shows that a strong $n$-partite tournament may not contain strong tournament with $n$ vertices. So Theorem 6 is not a trivial generalization of Corollary 10.

Now we would like to give the following related problems.
Problem 12. Are there examples of strong c-partite tournaments which are not tournaments with exactly $c-k+1(k-3)$-reducible $k$-cycles for $4 \cdot k \cdot c$ ?

The weak form of Problem 12 is also unsolved. We refer the reader who is interested in this problem to [7].

Problem 13 (Volkmann [7]). Are there examples of strong c-partite tournaments which are not tournaments with exactly $c-k+1$ cycles of length $k$ for $4 \cdot k \cdot c$ ?

In [8], Yao proved that for $T_{n}$ a strong tournament of order $n$, if there is an integer $k(3<k<n)$ such that $T_{n}$ contains exactly $n-k+1 k$-cycles, then $T_{n} \cong M_{n}$.

Problem 14. How to characterize extremal strong c-partite tournaments containing minimum number of cycles?

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