



ELSEVIER

Available online at www.sciencedirect.com

SCIENCE @ DIRECT®

European Journal of Combinatorics xx (xxxx) xxx–xxx

European Journal
of Combinatoricswww.elsevier.com/locate/ejc

The Ramsey numbers of trees versus W_6 or W_7 [☆]

Yaojun Chen, Yunqing Zhang, Kemin Zhang

Department of Mathematics, Nanjing University, Nanjing 210093, China

Received 6 September 2004; accepted 11 December 2004

1 Abstract

Let T_n denote a tree of order n and W_m a wheel of order $m + 1$. In this paper, we show the Ramsey numbers $R(T_n, W_6) = 2n - 1 + \mu$ for $n \geq 5$, where $\mu = 2$ if $T_n = S_n$, $\mu = 1$ if $T_n = S_n(1, 1)$ or $T_n = S_n(1, 2)$ and $n \equiv 0 \pmod{3}$, and $\mu = 0$ otherwise; $R(T_n, W_7) = 3n - 2$ for $n \geq 6$.

© 2005 Elsevier Ltd. All rights reserved.

Keywords: Ramsey number; Tree; Wheel

1. Introduction

All graphs considered in this paper are finite simple graphs without loops. For two given graphs G_1 and G_2 , the *Ramsey number* $R(G_1, G_2)$ is the smallest positive integer n such that for any graph G of order n , either G contains G_1 or \overline{G} contains G_2 , where \overline{G} is the complement of G . Let G be a graph and m be a positive integer. We use mG to denote m vertex disjoint copies of G . A *path* and a *cycle* of order n are denoted by P_n and C_n , respectively. A *star* S_n ($n \geq 3$) is a bipartite graph $K_{1,n-1}$. A *complete graph* of order n is denoted by K_n . A *wheel* $W_n = K_1 + C_n$ is a graph of order $n + 1$ vertices, where K_1 is called the *hub* of the wheel. $S_n(l, m)$ is a tree of order n obtained from $S_{n-l \times m}$ by subdividing each of l chosen edges m times. $S_n(l)$ is a tree of order n obtained from an S_l and an S_{n-l} by adding an edge joining the centers of them. $S_n[l]$ is a tree of order n obtained from an

[☆] This project was supported by NSFC under grant number 10201012.
E-mail address: yaojunc@nju.edu.cn (Y. Chen).

S_l and an S_{n-l} by adding an edge joining a vertex of degree one of S_l to the center of S_{n-l} . Define

$$\mathcal{T} = \{S_n \mid n \geq 5\} \cup \{S_n(1, 1) \mid n \geq 5\} \cup \{S_n(1, 2) \mid n \geq 6 \text{ and } n \equiv 0 \pmod{3}\}.$$

For a tree T , we define $L(T) = \{v \mid v \in V(T) \text{ and } d(v) = 1\}$. Let $V \subseteq L(T)$ and $|V| = k$. Write $T_V = T - V$. If $T_V \notin \mathcal{T}$, we call V a k -deletable set. If $k = 2$ and $|N(V)| = 2$, we call V a II-set. If $k = 3$ and $|N(V)| = 3$, we call V a III-set. If $k = 3$ and $|N(V)| = 2$, we call V a IV-set. If V is a II-set and $T_V \notin \mathcal{T}$, we call V a II-deletable set. Similarly, we can define III-deletable and IV-deletable sets. Terminology and notations not defined here can be found in [2].

In [1], Baskoro et al. obtain the following.

Theorem 1 ([1]). *Let T_n be a tree of order n other than S_n . Then $R(T_n, W_4) = 2n - 1$ for $n \geq 3$; $R(T_n, W_5) = 3n - 2$ for $n \geq 4$.*

Motivated by Theorem 1, Baskoro et al. [1] pose the following.

Conjecture 1. *Let T_n be a tree of order n other than S_n and $n \geq m - 1$. Then $R(T_n, W_m) = 2n - 1$ for even $m \geq 6$; $R(T_n, W_m) = 3n - 2$ for odd $m \geq 7$.*

In [3], we show Conjecture 1 holds for $T_n = P_n$.

Theorem 2 ([3]). *$R(P_n, W_m) = 3n - 2$ for m odd and $n \geq m - 1 \geq 2$; $R(P_n, W_m) = 2n - 1$ for m even and $n \geq m - 1 \geq 3$.*

In [4], we obtain the following.

Theorem 3 ([4]). *$R(S_n, W_6) = 2n + 1$ for $n \geq 3$; $R(S_n, W_m) = 3n - 2$ for m odd and $n \geq m - 1 \geq 2$.*

Using Theorem 3, we consider $R(T_n, W_6)$ for $\Delta(T_n) \geq n - 3$ in [5] and the following are established.

Theorem 4 ([5]). *$R(S_n(1, 1), W_6) = 2n$ for $n \geq 4$.*

Theorem 5 ([5]). *$R(S_n(1, 2), W_6) = 2n$ for $n \geq 6$ and $n \equiv 0 \pmod{3}$.*

Theorem 6 ([5]). *$R(S_n(3), W_6) = R(S_n(2, 1), W_6) = 2n - 1$ for $n \geq 6$; $R(S_n(1, 2), W_6) = 2n - 1$ for $n \geq 6$ and $n \not\equiv 0 \pmod{3}$.*

By Theorems 4 and 5, we can see that Conjecture 1 is not true when $m = 6$. In fact, as pointed out in [5], for even m , $R(T_n, W_m)$ is a function related to both n and m . However, we believe that $R(T_n, W_6) = 2n - 1$ for $T_n \notin \mathcal{T}$.

In [6], we evaluate $R(T_n, W_6)$ for $5 \leq n \leq 8$ and get the following.

Theorem 7 ([6]). *Let $T_n \notin \mathcal{T}$ be a tree of order n and $5 \leq n \leq 8$, then $R(T_n, W_6) = 2n - 1$.*

In [7], we consider $R(T_n, W_6)$ for T_n without certain deletable sets and establish the following.

Theorem 8 ([7]). Let $T \notin \mathcal{T}$ be a tree of order $n \geq 9$. If T contains no II-deletable set, or $|L(T)| \geq 3$ and T contains neither III-deletable set nor IV-deletable set, or $|L(T)| \geq 4$ and T contains no IV-deletable set, then $R(T, W_6) = 2n - 1$.

In this paper, we will determine $R(T_n, W_6)$ for all $T_n \notin \mathcal{T}$ and $n \geq 5$. On the other hand, we will consider the conjecture in the case where m is odd. As a special case, this paper will determine $R(T_n, W_7)$.

Let T_n be a tree of order n . The main results of this paper are the following.

Theorem 9. $R(T_n, W_6) = 2n - 1 + \mu$ for $n \geq 5$, where $\mu = 2$ if $T_n = S_n$, $\mu = 1$ if $T_n = S_n(1, 1)$ or $T_n = S_n(1, 2)$ and $n \equiv 0 \pmod{3}$, and $\mu = 0$ otherwise.

Theorem 10. $R(T_n, W_7) = 3n - 2$ for $n \geq 6$.

By **Theorem 10**, we can see that **Conjecture 1** holds for $m = 7$. For odd $m \geq 9$, the conjecture is still alive. Although the conjecture is not true for even m in general, we believe it holds for odd m .

2. Some lemmas

In order to prove the main results of this paper, we need the following lemmas.

Lemma 1 ([5]). Let G be a graph of order $2n - 1 \geq 7$ and (U, V) a partition of $V(G)$ with $|U| \geq 3$ and $|V| \geq 4$. Suppose $u_i \in U$ and $N_V(u_i) = \emptyset$, $1 \leq i \leq 3$. If \overline{G} contains no W_6 , then $\delta(G[V]) \geq |V| - 3$.

Lemma 2 ([7]). Let G be a graph of order $2n - 1$. If $\alpha(G) \leq 2$, then G contains all trees of order n .

Ore showed in [8] that if a graph on n vertices in which the degree sum of any two nonadjacent vertices is at least $n + 1$, then G is Hamilton-connected. From this result we can get easily the following.

Lemma 3. Let G be a graph of order n . If $\delta(G) \geq n/2 + 1$, then G is Hamilton-connected.

Lemma 4. Let G be a graph of order $n \geq 9$. If $\alpha(G) \leq 2$ and $\delta(G) \geq n - 3$, then G contains all trees T of order n with $|L(T)| = 3$.

Proof. If $\alpha(G) = 1$, then it holds trivially. Hence we may assume $\alpha(G) = 2$. Let T be a given tree of order n with $|L(T)| = 3$. Obviously, $\Delta(T) = 3$ and T has only one vertex of degree 3. Let $v \in V(G)$ and $G_0 = G - v$. Since $n \geq 9$ and $\delta(G) \geq n - 3$, we have $\delta(G_0) \geq (n - 3) - 1 \geq (n - 1)/2 + 1 = |G_0|/2 + 1$ which implies G_0 is Hamilton-connected by **Lemma 3**. If $d(v) = n - 3$, we assume $v_1, v_2 \notin N(v)$. Noting that $\alpha(G) = 2$, we have $v_1 v_2 \in E(G)$ and hence G_0 contains a Hamilton cycle $C = v_1 v_2 \cdots v_{n-1}$ such that $v_i \in N(v)$ for $3 \leq i \leq n - 1$. In this case, it is easy to see G contains T . If $d(v) \geq n - 2$, then since G_0 contains a Hamilton cycle, it is not difficult to see G contains T . \square

The following lemma is well known and can be found in many graph theory textbooks, see for instance [2].

Lemma 5. A bipartite graph G with bipartition (U, V) contains a matching saturating U if and only if $|N(S)| \geq |S|$ for every $S \subseteq U$.

3. Proof of Theorem 9

Proof of Theorem 9. Let T be a given tree of order $n \geq 5$. If $|L(T)| = 2$, then $T = P_n$ and hence Theorem 9 holds by Theorem 2. If $T \in \mathcal{T}$, then Theorem 9 holds by Theorems 3–5. Thus we may assume $|L(T)| \geq 3$ and $T \notin \mathcal{T}$.

We use induction on n . If $5 \leq n \leq 8$, then Theorem 9 holds by Theorem 7. In the following proof, we assume $n \geq 9$ and Theorem 9 holds for small values of n .

Let G be a graph of order $2n - 1$. If \overline{G} contains no W_6 , then $\alpha(G) \leq 6$. Let I be a maximum independent set of G . By Lemma 2, we may assume $|I| \geq 3$. Let $I = \{v_1, v_2, \dots, v_k\}$, where $3 \leq k \leq 6$.

Suppose to the contrary G contains no T . We now consider the following two cases.

Case 1. $k = 3$.

In order to prove Case 1, we need the following three claims.

Claim 1. G contains an induced subgraph $K_1 \cup K_2 \cup K_3$.

Proof. Since G contains no T , by Theorem 8 we may assume T contains a II-deletable set U_0 . By induction hypothesis, $G - I$ contains $T_{U_0} = T - U_0$. Let $N_T(U_0) = U$. If $|N_I(U)| \geq 2$, then G contains T , a contradiction. Hence $|N_I(U)| = 1$. This implies G has an induced subgraph $2K_1 \cup K_2$. Assume, without loss of generality, that $G[I_1] = 2K_1 \cup K_2$ with $I_1 = I \cup \{v_4\}$ and $v_3v_4 \in E(G)$. By induction hypothesis, $G - I_1$ contains T_{U_0} . If $|N_{I_1}(U)| \geq 2$, then G contains T , a contradiction. Hence $|N_{I_1}(U)| = 1$. Thus, noting that $k = 3$, we may assume $N_{I_1}(U) = \{v_2\}$. Let $I_2 = I_1 \cup U$. Since $k = 3$, it is easy to see that $G[I_2] = K_1 \cup K_2 \cup K_3$. \square

In the following, we let $G_0 = K_1 \cup K_2 \cup K_3$ with $V(G_0) = X = \{x_1, x_2, \dots, x_6\}$ and $E(G_0) = \{x_2x_3, x_4x_5, x_4x_6, x_5x_6\}$ be an induced subgraph of G .

Claim 2. $|L(T)| \geq 4$.

Proof. Let $L(T) = U_0$. If $|L(T)| = 3$, then by Theorem 8 we may assume U_0 is a IV-deletable set or III-deletable set. Let $N_T(U_0) = U$.

If $|U| = 2$, we assume $U = \{u_1, u_2\}$. In this case, it is easy to see $T_{U_0} = P_{n-3}$ and either $d_T(u_1) = 3$ or $d_T(u_2) = 3$. By Theorem 2, $G - I$ contains a P_{n-2} . Assume $P_{n-2} = p_1p_2 \cdots p_{n-2}$ to be a path in $G - I$. If $N_I(p_1) \cap N_I(p_{n-2}) \neq \emptyset$, then G contains a cycle C of length $n - 1$. Let $V = V(G) - V(C)$, then $d_V(v) = 0$ for any $v \in V(C)$ since otherwise G contains T . Thus we have $\alpha(G[V]) \leq 2$ since $\alpha(G) = 3$ and $\delta(G[V]) \geq n - 3$ by Lemma 1 and hence $G[V]$ contains T by Lemma 4. Thus we may assume $N_I(p_1) \cap N_I(p_{n-2}) = \emptyset$. In this case, if $d_I(p_1) \geq 2$ or $d_I(p_{n-2}) \geq 2$, then G contains T and hence we may assume $N_I(p_1) = \{v_1\}$ and $N_I(p_{n-2}) = \{v_2\}$. Let $V_0 = V(G) - I - P_{n-2}$, then $|V_0| = n - 2$. If $d_{V_0}(v_1) \geq 2$ or $d_{V_0}(v_2) \geq 2$, then G contains T . Thus, since $|V_0| = n - 2 \geq 7$, there are three vertices $w_1, w_2, w_3 \in V_0$ such that $v_iw_j \notin E(G)$ for $i = 1, 2$ and $j = 1, 2, 3$. If $N(p_2) \cap (V_0 \cup \{v_1\}) \neq \emptyset$ or

1 $N_{V_0}(p_{n-2}) \neq \emptyset$, then G contains T . Hence we have $N(p_2) \cap (V_0 \cup \{v_1\}) = N_{V_0}(p_{n-2})$
 2 $= \emptyset$. Thus, $\overline{G}[v_1, p_2, v_2, p_{n-2}, w_1, w_2, w_3]$ contains a W_6 with the hub v_1 , a contradiction.

3 If $|U| = 3$, we assume $U = \{u_1, u_2, u_3\}$. By induction hypothesis, $G - X$ contains
 4 T_{U_0} . Assume $d_X(u_1) \leq d_X(u_2) \leq d_X(u_3)$. Since G contains no T , by Lemma 5, we have
 5 $d_X(u_1) \leq d_X(u_2) \leq 2$. If $d_X(u_1) = 1$, then since $k = 3$, we have $N_X(u_1) = \{x_1\}$.
 6 If $x_1u_2 \in E(G)$, then since $d_X(u_2) \leq 2$, by the symmetry of x_2 and x_3 , we may
 7 assume $x_2u_2 \notin E(G)$. Thus $\overline{G}[x_2, u_1, u_2, x_1, x_4, x_5, x_6]$ contains a W_6 with the hub
 8 x_2 , a contradiction. If $x_1u_2 \notin E(G)$, then since $d_X(u_2) \leq 2$ and $k = 3$, we must
 9 have $N_I(u_2) = \{x_2, x_3\}$. In this case, $\overline{G}[x_1, u_2, x_2, x_3, x_4, x_5, x_6]$ contains a W_6 with
 10 the hub x_1 , again a contradiction. Thus we have $d_X(u_1) = 2$. If $|N_X(U)| \geq 3$, then by
 11 Lemma 5, G contains T , a contradiction. Hence we have $|N_X(U)| = 2$ which implies
 12 $N_X(u_1) = N_X(u_2) = N_X(u_3)$. Let $W = V(G) - N_X(U) - V(T_{U_0})$, then $|W| = n$.
 13 Obviously, $d_W(u_i) = 0$ for $i = 1, 2, 3$ since otherwise G contains T . This implies
 14 $\alpha(G[W]) \leq 2$ since $k = 3$. And then $\delta(G[W]) \geq n - 3$ by Lemma 1. Hence $G[W]$
 15 contains T by Lemma 4, a contradiction. \square

16 **Claim 3.** G contains no induced subgraph $3K_2$.

17 **Proof.** Since G contains no T , by Theorem 8 we may assume T contains a III-deletable
 18 set or IV-deletable set U_0 and $N_T(U_0) = U$. Suppose to the contrary G contains an
 19 induced subgraph $G_1 = 3K_2$ with $V(G_1) = Y = \{y_i \mid 1 \leq i \leq 6\}$ and $E(G_1) =$
 20 $\{y_1y_2, y_3y_4, y_5y_6\}$. By induction hypothesis, $G - Y$ contains T_{U_0} . Since $k = 3$, it is easy
 21 to see $d_Y(u) \geq 2$ for any $u \in U$. If $|N_Y(U)| \geq 3$ and $|U| = 2$, then it is easy to see
 22 G contains T , a contradiction. If $|N_Y(U)| \geq 3$ and $|U| = 3$, then G contains T by
 23 Lemma 5, a contradiction. Thus we have $|N_Y(U)| = 2$. Since $k = 3$, we may assume
 24 $N_Y(U) = \{y_5, y_6\}$. In this case, G contains an induced subgraph $G_2 = 2K_2 \cup K_4$. Let
 25 $V(G_2) = Z = \{z_i \mid 1 \leq i \leq 8\}$ and $E(G_2) = \{z_1z_2, z_3z_4\} \cup \{z_iz_j \mid 5 \leq i < j \leq 8\}$. By
 26 Claim 2, we have $|L(T)| \geq 4$. By Theorem 8 we may assume T contains a IV-deletable
 27 set U_1 . If $d_Z(u) \geq 4$ for any $u \in N_T(U_1)$, then G contains T , a contradiction. Hence there
 28 is some vertex $u_0 \in N_T(U_1)$ such that $d_Z(u_0) \leq 3$. Set $V = \{z_5, z_6, z_7, z_8\}$. Since $k = 3$,
 29 we have $d_V(u_0) \leq 1$. Hence we may assume $N_Z(u_0) \cap \{z_5, z_6, z_7\} = \emptyset$. Since $d_Z(u_0) \leq 3$,
 30 we may assume $z_1 \notin N_Z(u_0)$. Thus $\overline{G}[z_1, u_0, z_3, z_4, z_5, z_6, z_7]$ contains a W_6 with the hub
 31 z_1 , a contradiction. \square

32 In the following we prove Case 1.

33 By Theorem 8, T contains a III-deletable set or a IV-deletable set U_0 . Let $N_T(U_0) = U$.
 34 By induction hypothesis, $G - X$ contains T_{U_0} . If there is some vertex $u \in U$ such
 35 that $d_X(u) = 1$, then since $k = 3$, we have $N_X(u) = \{x_1\}$ and thus G contains
 36 an induced subgraph $3K_2$ which contradicts Claim 3. Hence we have $d_X(u) \geq 2$ for
 37 any $u \in U$. If $|N_X(U)| \geq 3$ and $|U| = 2$, then G contains T , a contradiction. If
 38 $|N_X(U)| \geq 3$ and $|U| = 3$, then G contains T by Lemma 5, a contradiction. Hence we
 39 have $|N_X(U)| = 2$ which implies $N_X(u) = N_X(U)$ for each $u \in U$. If $x_1 \in N_X(U)$,
 40 then by the symmetry of x_2 and x_3 and Claim 3, we may assume $x_3 \in N_X(U)$ and
 41 hence $\overline{G}[x_2, u_1, u_2, x_1, x_4, x_5, x_6]$ contains a W_6 with the hub x_2 , where $u_1, u_2 \in U$, a
 42 contradiction. If $x_1 \notin N_X(U)$, then since $k = 3$, we have $N_X(U) = \{x_2, x_3\}$ which implies
 43 $\overline{G}[x_1, u_1, x_2, x_3, x_4, x_5, x_6]$ contains a W_6 with the hub x_1 , where $u_1 \in U$, a contradiction.

Case 2. $k \geq 4$.

If $k = 4$, then by [Theorem 8](#) we may assume T contains a II-deletable set U_0 . By induction hypothesis, $G - I$ contains T_{U_0} . Let $N_T(U_0) = U$. If $|N_I(U)| \geq 2$, then G contains T , a contradiction. Thus we have $|N_I(U)| = 1$ which implies G contains an induced subgraph $3K_1 \cup K_3$. Let $G' = 3K_1 \cup K_3$ with $V(G') = W = \{w_i \mid 1 \leq i \leq 6\}$ and $E(G') = \{w_4w_5, w_4w_6, w_5w_6\}$. By [Theorem 8](#) we may assume T contains a III-deletable set U_1 . Let $N_T(U_1) = U_2$. By induction hypothesis, $G - W$ contains T_{U_1} . If $d_W(u) \geq 3$ for each $u \in U_2$, then G contains T , a contradiction. Hence there is some vertex $u_0 \in U_2$ such that $d_W(u_0) \leq 2$. Since $k = 4$, we have $|N(u_0) \cap \{w_4, w_5, w_6\}| \leq 1$. Since $d_W(u_0) \leq 2$, we may assume $w_1 \notin N(u_0)$. Thus $\overline{G}[w_1, w_2, w_3, u_0, w_4, w_5, w_6]$ contains a W_6 with the hub w_1 , a contradiction.

Let now $k = 5, 6$. By [Theorem 8](#) we may assume T contains a 3-deletable set U_0 . Let $N_T(U_0) = U$. By induction hypothesis, $G - I$ contains T_{U_0} . If $d_I(u) \geq 3$ for each $u \in U$, then G contains T , a contradiction. Hence there is some vertex $u \in U$ such that $d_I(u) \leq 2$. Thus, if $k = 5$, then G contains an induced subgraph $3K_1 \cup P_3$ or $4K_1 \cup K_2$. By an analogous argument of $k = 4$, we can get a contradiction. If $k = 6$, then $\overline{G}[I \cup \{u\}]$ contains a W_6 , a contradiction.

From the proof above, we have $R(T, W_6) \leq 2n - 1$ for $T \notin \mathcal{T}$. On the other hand, the graph $2K_{n-1}$ shows $R(T, W_6) \geq 2n - 1$ for any tree T of order n and hence $R(T, W_6) = 2n - 1$ for $T \notin \mathcal{T}$. Thus the proof of [Theorem 9](#) is completed. \square

4. Proof of Theorem 10

Proof of Theorem 10. Let G be a graph of order $3n - 2$ and T a given tree of order n . Suppose \overline{G} contains no W_7 .

Claim 4. *If G contains no T , then $\delta(G) = n - 2$.*

Proof. By [Theorem 3](#), we may assume $T \neq S_n$. Let $d(v) = \delta(G)$ and $V = V(G) - N[v]$. If $\delta(G) \leq n - 3$, then $|V| \geq 2n$. Since G contains no T , by [Theorem 9](#), $\overline{G}[V]$ contains a W_6 and hence $\overline{G}[V]$ contains a C_7 which implies \overline{G} contains a W_7 with the hub v , a contradiction. Hence we have $\delta(G) \geq n - 2$. If $\delta(G) \geq n - 1$, then it is easy to see G contains all trees of order n . Thus we have $\delta(G) \leq n - 2$ and hence $\delta(G) = n - 2$. \square

By [Theorem 3](#), G contains a tree $T_* = S_n$. Let $V(T_*) = \{v_0, v_1, \dots, v_{n-1}\}$ and $E(T_*) = \{v_0v_i \mid 1 \leq i \leq n - 1\}$. If G contains no $S_n(1, 1)$, then by [Claim 4](#), we have $d(v_1) \geq n - 2 \geq 4$ and hence there is some vertex $w \in V(G)$ such that $w \neq v_0$ and $w \in N(v_1)$ which implies G contains an $S_n(1, 1)$, a contradiction. Assume $T_{**} = S_n(1, 1)$ with $V(T_{**}) = \{u_0, \dots, u_{n-1}\}$ and $E(T_{**}) = \{u_0u_i \mid 1 \leq i \leq n - 2\} \cup \{u_1u_{n-1}\}$. If G contains no $S_n(1, 2)$, then by [Claim 4](#), we have $d(u_{n-1}) \geq n - 2 \geq 4$ and hence there is some vertex $w \in V(G)$ such that $w \neq u_0, u_1$ and $w \in N(u_{n-1})$ which implies G contains an $S_n(1, 2)$, a contradiction. Thus we may assume $T \neq S_n, S_n(1, 1)$ and $S_n(1, 2)$. Assume $d(v) = \delta(G)$ and $V = V(G) - N[v]$. If G contains no T , then by [Claim 4](#), we have $|V| = 2n - 1$. By [Theorem 9](#), $\overline{G}[V]$ contains a W_6 and hence $\overline{G}[V]$ contains a C_7 which implies \overline{G} contains a W_7 with the hub v , a contradiction. Thus we have

1 $R(T, W_7) \leq 3n - 2$. On the other hand, the graph $3K_{n-1}$ shows $R(T, W_7) \geq 3n - 2$
2 and hence $R(T, W_7) = 3n - 2$, that is, $R(T_n, W_7) = 3n - 2$. The proof of Theorem 10 is
3 completed. \square

4 Acknowledgements

5 The first author was supported partially by Nanjing University Talent Development
6 Foundation. The third author was supported partially by NSFZJ.

7 References

- 8 [1] E.T. Baskoro, Surahmat, S.M. Nababan, M. Miller, On Ramsey numbers for trees versus wheels of five or six
9 vertices, *Graphs and Combinatorics* 18 (2002) 717–721.
10 [2] J.A. Bondy, U.S.R. Murty, *Graph Theory with Applications*, North-Holland, Amsterdam, 1976.
11 [3] Y.J. Chen, Y.Q. Zhang, K.M. Zhang, The Ramsey numbers of paths versus wheels, *Discrete Mathematics* 290
12 (2005) 85–87.
13 [4] Y.J. Chen, Y.Q. Zhang, K.M. Zhang, The Ramsey numbers of stars versus wheels, *European Journal of*
14 *Combinatorics* 25 (2004) 1067–1075.
15 [5] Y.J. Chen, Y.Q. Zhang, K.M. Zhang, The Ramsey numbers $R(T_n, W_6)$ for $\Delta(T_n) \geq n - 3$, *Applied*
16 *Mathematics Letters* 17 (2004) 281–285.
17 [6] Y.J. Chen, Y.Q. Zhang, K.M. Zhang, The Ramsey numbers $R(T_n, W_6)$ for small n , *Utilitas Mathematica*
18 (in press).
19 [7] Y.J. Chen, Y.Q. Zhang, K.M. Zhang, The Ramsey numbers $R(T_n, W_6)$ for T_n without certain deletable sets,
20 *Journal of Systems Science and Complexity* 18 (2005) 95–101.
21 [8] O. Ore, Arc coverings of graphs, *Annali di Matematica Pura ed Applicata* 55 (1961) 315–321.