The Minimal Solutions of Boolean Matrix-equation $A^k = J$

MIAO Zheng ke *) (苗正科) and ZHANG Ke min (张克民) (Department of Mathematics, Nanjing University, Nanjing, 210093)

Abstract: Let A be a primitive Boolean matrix. (A) is the least number k such that $A^k = J$. (A) is the number of 1-entries in A. In this paper, the parameter $N(k, n) = \min\{(A) \mid A^T = A, \operatorname{trace}(A) = 0, (A) = k\}$ is considered. Furthermore, we describe the set $EG(k, n) = \{G(A) \mid (A) = N(k, n), A^T = A, \operatorname{trace}(A) = 0, (A) = k\}$ and obtain a characterization of the minimal solutions with zero trace of the Boolean matrix equation $A^k = J$.

Key words: primitive, exponent, norm

1991 MR subject classification: 05C20, 05C50, 15A33, 15A24

CLC number: O157.5,O151.21

Document code: A

Article ID: 1000-1778(2000)02-0155-05

Let $B = \{0, 1\}$ be the usual binary Boolean algebra. The matrices over B are called Boolean matrices. An $n \times n$ Boolean matrix A is called a primitive matrix if there exists a positive integer k such that $A^k = J$ (where J is the universal matrix). The least such k is called the exponent of A, denoted by (A).

In projective plane theory, although a lot of results are obtained on Boolean matrix equation, it is still a famous open problem to find the square roots of a Boolean matrix.

Let A be an $n \times n$ Boolean matrix. Define the norm of A, denoted by (A), to be the number of 1-entries in A. Clearly, satisfies the norm axioms. As you know, a special Boolean matrix-equation $A^k = J$ has solutions. In general, it is very difficult to solve this equation. So, the parameter

$$N(k, n) = \min\{ (A) / A^T = A, \operatorname{trace}(A) = 0, (A) = k \}$$

must be considered. And we need the following concepts and propositions. The associated graph of an $n \times n$ symmetric Boolean matrix $A = (a_{ij})$, denoted by G(A), is the graph

Received date: Jan. 19, 1998.

Foundation item: The NSF (19871040) of China and NSF (BK97041105) of Jiangsu Province.

^{*)} Permanent address: Department of Mathematics, Xuzhou Normal University, Xuzhou, 221009.

with vertex set $V = \{1, 2, ..., n\}$ such that there is an edge arc between i and j in D(A) if and only if $a_{ij} = a_{ji} = 1$. A graph G is primitive if there exists an integer k > 0 such that for all pairs of vertices i, j V(G) (not necessary distinct), there is a walk from i to j with length k. The least such k is called the exponent of G, denoted by G(G). Clearly, a symmetric Boolean matrix G(G) is primitive and G(G).

Let G be a primitive graph with order n. For any $i, j \in V(G)$, the local exponent from i to j, denoted by (i, j), is the least integer k such that there exists a walk of length m from i to j for all $m \in k$. It is obvious that $(G) = \max_{i \in V(G)} (i, j)$.

Proposition^[1] The exponent set of symmetric primitive (0,1)-matrices with zero trace is $\mathcal{E}_n = \{2,3,...,2n-4\}$ - Y, where Y is the set of all odd numbers in $\{n-2,n-1,...,2n-5\}$.

In this paper, we describe the set

 $EG(k,n) = \{ G(A) \mid (A) = N(k,n), A^T = A, \operatorname{trace}(A) = 0, (A) = k \}$ and give a characterization of the minimal symmetric solutions with zero trace of $A^k = J$. Other terms and notations not defined here, we refer the reader to [2].

According to the definition of N(k, n) and Proposition 3, $k \in \mathcal{E}_n$.

Theorem 1 $N(2, n) = 2 \left[\frac{3n-3}{2} \right]$ for n = 3. Moreover, G = EG(2, n) is unique in the sense of permutation similarity.

Proof Since $(G_1) = (G_2) = 2$ for n = 3 (see Figure 1), we have $N(2, n) = 2 \left[\frac{3n-3}{2} \right]$.

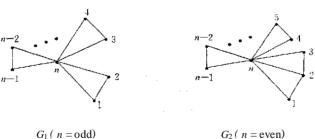


Figure 1

If $A^T = A$, $A^2 = J$ and trace (A) = 0, then there is a walk with length 2 from one vertex to another in G(A). So any vertex of G(A) is on some 3-cycle. Otherwise, there would exist a vertex u = V(G(A)) such that u is a pendant vertex or u is on an s-cycle (s = 4). Let v be a neighbouring vertex of u. Then (v, u) > 2, a contradiction. Since $A^2 = J$, for any u, v = V(G(A)) there exists 3-cycles C_1 , C_2 such that $u = C_1$, $v = C_2$ and $C_1 = C_2$ \emptyset . For any 3-cycle C, let

 $t = \begin{cases} u : \text{ There exists a 3-cycle } C \text{ containing } u \\ \text{such that } C \text{ is exactly one vertex} \end{cases}$

Thus we have

$$(G(A)) \quad 3 + \left[\frac{3t}{2}\right] + 2(n - 3 - t) = \left[\frac{2n - 3 - t}{2}\right] \\ \left[2n - 3 - \frac{n - 3}{2}\right] = \left[\frac{3n - 3}{2}\right]. \tag{1}$$

So (A) $2\left[\frac{3n-3}{2}\right]$. Hence $N(2, n) = 2\left[\frac{3t}{2}\right]$.

If G = EG(2, n), then (1) is an equality. So $G \cong G_1$ when n is odd and $G \cong G_2$ when n is even.

Theorem 2 $N(3,n) = 2\left[\frac{3n-4}{2}\right]$ for n=6. Further, G=EG(3,n) is unique in the sense of permutation similarity.

Proof Since $(G_3) = (G_4) = 3$ for n = 6 (see Figure 2), we have $N(3, n) = 2 \left[\frac{3n-4}{2} \right]$.

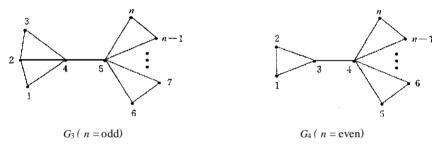


Figure 2

Let $A^T = A$, trace (A) = 0 and (A) = 3. Then for any u = V(G(A)) there is a 3-cycle containing u. Otherwise, (u, u) > 3. By (A) = 3, the diameter of G(A) is less then 4 and greater then 1.

Case 1 The diameter of G(A) is equal to 3.

We claim that G(A) has a subgraph H (see Figure 3).



Figure 3 H

We take any two 3-cycles C_1 , C_2 , and let $d=d_G(C_1,C_2)-3$. Thus there exists a path with length d. For d=3 (similarly for d=2), let P=xuvy, $x-C_1$ and $y-C_2$. If there is a 3-cycle C_3 containing v such that $C_2-C_3=\emptyset$, then the claim holds. Hence C_2-C_3 . \emptyset . Now, we consider a 3-cycle C_4 with $u-C_4$. Thus we have $C_1-C_4-\emptyset$ and $C_3-C_4-\emptyset$. Hence $d_G(C_1,C_2)-2$, a contradiction. So G(A) contains a subgraph H.

Case 2 The diameter of G(A) is equal to 2.

If there exist two 3-cycles C_1 , C_2 such that u C_1 , v C_2 and C_1 C_2 \emptyset for any u, v V(G(A)), then (A) = 2. This is a contradiction. If there exist two vertices u, v and two 3-cycles C_1 , C_2 such that u C_1 , v C_2 and C_1 $C_2 = \emptyset$, then we can prove that

G(A) has a subgraph H as in the proof of the case 1. Let

$$t = \begin{cases} u : \text{ There exists a 3-cycle } C \text{ containing } u \\ \text{such that } C & H \text{ is exactly one vertex} \end{cases} \end{cases}$$

Thus we have

$$(G(A)) 7 + \left[\frac{3t}{2}\right] + 2(n - 6 - t) = \left[\frac{2n - 5 - t}{2}\right]$$

$$\left[2n - 5 - \frac{n - 6}{2}\right] = \left[\frac{3n - 4}{2}\right]. (2)$$

So
$$(A)$$
 $2\left[\frac{3n-4}{2}\right]$. Hence $N(3,n) = 2\left[\frac{3n-4}{2}\right]$.

If G = EG(3, n), then (2) is an equality. So $G \cong G_3$ when n is odd and $G \cong G_4$ when n is even.

Theorem 3 N(2q, n) = 2n for 2 - q - n - 2.

Proof Since $(G_5) = 2q$ for 2 = q = n - 2 (see Figure 4), we have N(2q, n) = 2n for 2 = q = n - 2.

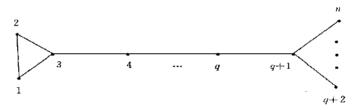


Figure 4

If A is a symmetric primitive matrix with zero trace, then (A) 2 n. Hence N (2q, n) = 2n.

Lemma Let G be a primitive undirected graph with exactly one cycle C. Then (G) is even.

Proof Let G_1 , G_2 , ..., G_r be the components of G - E(C), and C the length of C. Let $t_i = \max_{v_i \ V(C_i)} \min_{u \ V(C)} d(u, v_i)$, $t = \max\{t_1, t_2, ..., t_r\}$.

It is obvious that there exists a vertex u_0 such that the distance from u_0 to C is t. Hence $(u_0, u_0) = 2t + c - 1$.

If u, v $V(G_i)$, then

$$(u, v)$$
 $2t_i + c - 1$ $2t + c - 1$.

If $u = V(G_i)$, $v = V(G_i)(j = i)$, then

$$(u, v)$$
 $t_i + t_i + c - 1$ $2t + c - 1$.

Therefore (G) = 2t + c - 1 is even.

Theorem 4
$$N(2q+1, n) = 2(n+1)$$
 for 2 $q = \left[\frac{n-4}{2}\right]$

Proof If (A) is odd, then G(A) has at least two odd cycles by the above Lemma. So N(2q+1, n) = 2(n+1).

Since
$$(G_6) = 2q + 1$$
 for $2 q \left[\frac{n-4}{2}\right]$ (see Figure 5), we have

$$N(2q+1, n) = 2(n+1)$$
 for $2 q $\left[\frac{n-4}{2}\right]$$

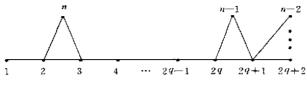


Figure 5

We assume that G is a primitive graph of order n with one cycle C exactly. Clearly, A(G) = 2n. Let $m = \max\{d(u, C) : u \mid V(G)\}$ and c be the length of C. We denote G by T(c, m). By the above Lemma (G) = 2m + c - 1 if c is odd. So $EG(2q, n) \supset \{T(c, m) : 2m + c - 1 = 2q\}$. On the other hand, if G = EG(2q, n), then G is a graph with no loop and (G) = 2q, (A(G)) = 2n. So G is a connected undirected graph with exactly one odd cycle C. Let c(>1) be the length of C, and m the maximum length of the path from a vertex to C in G. Thus (G) = 2m + c - 1 = 2q. So $G = T(c, m) = \{T(c, m) : 2m + c - 1 = 2q\}$. So we have M ain Result A is the minimal symmetric solution with zero trace of $A^k = J$ if and only if G(A) = (A) =

References

- [1] Liu Bolian, Mc Kay, B. D., Wormald, N. and Zhang Kemin, The exponent set of symmetric primitive (0,1)-matrices with zero trace, *Linear Algebra Appl.*, **136**(1990), 107 117.
- [2] Bondy J. A. and Murty , U. S. R., Graph Theory with Applications, Macmillan, London, 1976.