

A SUFFICIENT CONDITION FOR A GRAPH TO BE HAMILTONIAN*

Song ZengMin

(Southeast University)

Zhang Kemin

(Nanjing University)

Hamilton 图的一个充分条件

宋增民

(东南大学)

张克民

(南京大学)

摘 要

设 G 是 $n (\geq 3)$ 阶的 2-连通无向图. 若对任意的 $u, v \in V(G)$ 且 $d(u, v) = 2$ 均满足条件 $\text{Max}\{d(u), d(v)\} \geq \frac{n}{2}$ 或 $|N(u) \cup N(v)| \geq \frac{2}{3}(n-1)$, 则 G 是 hamiltonian. 在类似的条件下, 也讨论了图的 hamilton-connected 性质, 得到了相应的结果.

Abstract

In this paper we prove that a 2-connected graph G of order $n \geq 3$ is hamiltonian if for all distinct vertices u and v , $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq n/2$ or $|N(u) \cup N(v)| \geq (2n-2)/3$. We also demonstrate hamilton-connected property in graphs under similar conditions.

Introduction

This paper uses terms and notations of [2]. Throughout this paper G denotes an undirected 2-connected simple graph of order $n (\geq 3)$ with connectivity k and

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independence number α . Let L be a subset of $V(G)$, F a subgraph of G and v a vertex in G . Define $N_L(v) = \{u \in L | uv \in E(G)\}$, $N_L(F) = \bigcup_{v \in V(F)} N_L(v)$. Specifically, if $L = V(G)$, we simply write them as $N(v)$ and $N(F)$. If no ambiguity can arise we sometimes write F instead of $V(F)$.

The following results are the inspiration for the work in this paper.

Theorem A^[3] Let G be a 2-connected graph on n vertices. If for all distinct vertices u, v , $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq n/2$, then G is hamiltonian.

Theorem B^[4] Let G be a 2-connected graph of order n . If for all distinct vertices u, v , $d(u, v) = 2$ implies that $|N(u) \cup N(v)| \geq (2n - 1) / 3$, then G is hamiltonian.

In this paper, we shall prove a stronger result. Theorem A and B are corollaries of our result.

Main results

Theorem 1 Let G be a 2-connected graph of order $n \geq 3$. If for all distinct vertices u, v , $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq n/2$ or $|N(u) \cup N(v)| \geq (2n - 2) / 3$, then G is hamiltonian.

Proof It is trivial for $n \leq 4$, so we assume that $n \geq 5$. Let $A = \{u \in V(G) | d(u) \geq n/2\}$, $E' = \{xy | x, y \in A, xy \notin E(G)\}$ and $H = G + E'$. Then $H[A]$ is complete and there exists a cycle containing A in H . By Bondy and Chvatal's Closure Theorem^[1], G is hamiltonian if and only if H is hamiltonian. Thus, we only need to prove that H is hamiltonian. Let $C = v_1 v_2 \dots v_t v_1$ (simply, $12 \dots t1$) be a longest cycle containing A in H . If H is not hamiltonian, let B be any component of $H \setminus V(C)$. Let $N_C(B) = \{i_1, i_2, \dots, i_m\}$, $N^- = \{i_1 - 1, i_2 - 1, \dots, i_m - 1\}$ and $N^+ = \{i_1 + 1, i_2 + 1, \dots, i_m + 1\}$ where $i_1 < i_2 < \dots < i_m$, and where and later on $i \pm j$ is taken modulo t . Since $\chi(H) \geq \chi(G) = k$, we have

Assertion 1. $m \geq k \geq 2$

Let x_j be some vertex in B which is adjacent to i_j . It is possible that $x_i = x_j$ for $x \neq j$.

Assertion 2. For any j with $1 \leq j \leq m$, $d_H(x_j, i_j \pm 1) = 2$, and if $d_H(i_j \pm 1) < n/2$, then $d_G(x_j, i_j \pm 1) = 2$.

Assertion 3. For any j with $1 \leq j \leq m$, we have

$$(1) \quad d(i_j - 1) \leq n - |N(i_{j+1} - 1) \cup N(x_{j+1})| - \varepsilon,$$

$$(2) \quad d(i_j + 1) \leq n - |N(i_{j-1} + 1) \cup N(x_{j-1})| - \varepsilon,$$

where $\varepsilon = 0$, if $(i_j - 1)(i_j + 1) \in E(H)$, and $\varepsilon = 1$, if $(i_j - 1)(i_j + 1) \notin E(H)$.

Let $u \in N(i_{j+1} - 1) \cup N(x_{j+1})$. A bijection f is defined by:

$$f(u) = \begin{cases} u & \text{if } u \notin V(C). \\ i-1 & \text{if } u = i \in V(C) \text{ and } i_{j+1} \leq i \leq i_j - 1, \\ i+1 & \text{if } u = i \in V(C) \text{ and } i_j + 1 \leq i < i_{j+1} - 2, \\ i_j - 1 & \text{if } u = i_j \in V(C), \\ x_{i_{j+1}} & \text{if } u = i_{j+1} - 2 \in V(C) \text{ and } i_{j+1} - 2 \neq i_j. \end{cases}$$

Since C is a longest cycle in H , it is easy to check that $N(i_j - 1) \cap f(N(i_{j+1} - 1) \cup N(x_{i_{j+1}})) = \emptyset$ (For example, see [4]). So $d(i_j - 1) \leq n - |N(i_{j+1} - 1) \cup N(x_{i_{j+1}})|$. And note that when $(i_j - 1)(i_j + 1) \notin E(H)$, $x_{i_{j+1}} \notin f(N(i_{j+1} - 1) \cup N(x_{i_{j+1}}))$ if $i_{j+1} - 2 = i_j$ and $i_j + 1 \notin f(N(i_{j+1} - 1) \cup N(x_{i_{j+1}}))$ if $i_{j+1} - 2 > i_j$. Hence (1) is true. Similarly, (2) is true too.

Assertion 4. For any $u, v \in N^+$ or $u, v \in N^-$, $d(u) + d(v) < n$.

Let $u = i, v = j (i < j)$ and $i, j \in N^-$. At most one of $\{ik, (k-1)j\}$ belongs to $E(H)$, if $i < k < j$. So does $\{ik, (k+1)j\}$, if $j < k < i$. So $d_c(i) + d_c(j) \leq |C|$. On the other hand, $N(i) \cap N(j) \subseteq V(C)$. Hence $d(i) + d(j) = d(u) + d(v) < n$. Similarly, we have that $d(u) + d(v) < n$ for any $u, v \in N^+$.

Assertion 5. For any $u \in N^- \cup N^+$, $d(u) < n/2$.

If not, there is $u \in N^- \cup N^+$ such that $d(u) \geq n/2$, without loss of generality, we assume that $d(i_1 - 1) \geq n/2$. By assertion 4 for any $2 \leq j \leq m$, $d(i_j - 1) < n/2$. Hence, by assertion 2, $d_G(i_j - 1, x_j) = 2$. And then by the hypothesis of Theorem and assertion 1, $|N(i_2 - 1) \cup N(x_2)| \geq (2n - 2)/3$. Hence by assertion 3, $d(i_1 - 1) \leq (n + 2)/3 < n/2$, a contradiction.

By assertions 1, 2, 5 and the hypothesis of Theorem, we have:

Assertion 6. For any j with $1 \leq j \leq m$, $|N(i_j \pm 1) \cup N(x_j)| \geq (2n - 2)/3$.

If there exists an integer j , $1 \leq j \leq m$, with $(i_j - 1)(i_j + 1) \notin E(H)$, then $d_G((i_j - 1)(i_j + 1)) = 2$ and, by assertions 3, 6, $d(i_j - 1) \leq (n - 1)/3$, $d(i_j + 1) \leq (n - 1)/3$. this implies that $|N(i_j - 1) \cup N(i_j + 1)| \leq 2(n - 1)/3 - 1 = (2n - 5)/3$. This contradicts the hypothesis of the Theorem. Therefore,

Assertion 7. $(i_j - 1)(i_j + 1) \in E(H)$ ($1 \leq j \leq m$).

There exists a vertex h , $i_2 + 1 \leq h \leq i_3 - 1$ if $m \geq 3$ or $i_2 + 1 \leq h \leq i_1 - 1$ if $m = 2$ such that $h(i_2 - 1) \notin E(H)$, $i(i_2 - 1) \in E(H)$ for all $i_2 \leq i \leq h - 1$. This is true since $(i_2 - 1)(i_3 - 1) \notin E(H)$. or a cycle longer than C exists. Assertion 7 implies that $h \geq i_2 + 2$. Let $u \in N(i_1 + 1) \cup N(x_1)$. It is easy to prove that $u \notin \{i_2 + 1, i_2 + 2, \dots, h\}$. A bijection g is defined by:

$$g(u) = \begin{cases} u & \text{if } u \notin V(C), \\ i-1 & \text{if } u = i \in V(C) \text{ and } i_1 + 3 \leq i \leq i_2, \\ i+1 & \text{if } u = i \in V(C) \text{ and } h + 1 \leq i < i_1, \\ h & \text{if } u = i_1 + 2, \end{cases}$$

since C is a longest cycle in H , $g(N(i_1 + 1) \cup N(x_1)) \cap N(h) = \emptyset$. Note that $x_1 \notin g(N(i_1 + 1) \cup N(x_1)) \cup N(h)$. Thus by assertions 6, 7 and the hypothesis of Theorem, $d(h) \leq n - 2 - |N(i_1 + 1) \cup N(x_1)| \leq (n - 4) / 3$, if $i_2 h \notin E(H)$; $d(h) \leq n - 1 - |N(i_1 + 1) \cup N(x_1)| \leq (n - 1) / 3$ and $|N(i_2 - 1) \cap N(h)| \geq 2$, if $i_2 h \in E(H)$. On the other hand, by assertion 3, $d(i_2 - 1) \leq (n + 2) / 3$. Hence we have: (a) $\max\{d(i_2 - 1), d(h)\} < n / 2$; (b) $|N(i_2 - 1) \cup N(h)| \leq (n - 4) / 3 + (n + 2) / 3 - 1 = (2n - 5) / 3$ if $i_2 h \notin E(H)$; or $|N(i_2 - 1) \cup N(h)| \leq (n - 1) / 3 + (n + 2) / 3 - 2 = (2n - 5) / 3$ if $i_2 h \in E(H)$. This is contrary to the hypothesis of Theorem.

Consider the graph G_0 , which consists of three copies of K_r graphs G_1, G_2, G_3 and the set edges $\{x_1 x_2, x_2 x_3, x_3 x_1, y_1 y_2, y_2 y_3, y_3 y_1\}$, where $x_i, y_i \in V(G_i)$ for any $i, 1 \leq i \leq 3$. If $r \geq 3$, then it is easy to check that for any $u, v \in V(G_0)$ with $d(u, v) = 2$, $\max\{d(u), d(v)\} < n / 2$, $|N(u) \cup N(v)| \geq (2n - 3) / 3$ and G_0 is nonhamiltonian. So the conditions of Theorem 1 is the best possible in a sense.

We now consider hamilton-connected property. The graph G being 3-connected must be necessary, since there does not exist any $u-v$ hamiltonian path for any vertex cut $\{u, v\}$ in G .

Theorem 2 Let G be a 3-connected graph of order $n (\geq 3)$, and let u and v be distinct vertices of G . If $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq (n + 1) / 2$ or $|N(u) \cup N(v)| \geq 2n / 3$, then G is hamilton-connected.

Proof It is trivial for $n \leq 4$. So we assume that $n \geq 5$. Let $A = \{u \in V(G) | d(u) \geq (n + 1) / 2\}$, $E' = \{xy | x, y \in A, xy \notin E(G)\}$ and $H = G + E'$. Then $H[A]$ is complete. It is easy to prove that there exists a $u-v$ path containing A for any $u, v \in V(G)$ in H . By Bondy and Chvatal's Closure Theorem^[1], G is hamilton-connected if and only if H is hamilton-connected. Thus, we only need to prove that H is hamilton-connected. Suppose that H is not hamilton-connected. Then there exists a pair of vertices u and v such that no hamiltonian $u-v$ path exists in H . Consider a longest $u-v$ path P containing A , denoted $P = v_1 v_2 \dots v_t$ (simply, $1 2 \dots t$) in H , where $v_1 = u, v_t = v$. let B be any component of $H \setminus V(P)$. Let $N_p(B) = \{i_1, i_2, \dots, i_m\}, N^- = \{i_1 - 1, i_2 - 1, \dots, i_m - 1\}$ and $N^+ = \{i_1 + 1, i_2 + 1, \dots, i_m + 1\}$, where $i_1 < i_2 < \dots < i_m$. We can use an analogous arguments of Theorem 1 and have several assertions as follows:

Assertion 1. $m \geq k \geq 3$.

Let x_j be some vertex in B which is adjacent to i_j . It is possible that $x_i = x_j$ for $i \neq j$.

Assertion 2. For any j with $1 \leq j \leq m$, $d_H(x_j, i_j \pm 1) = 2$, and if $d_H(i_j \pm 1) \leq n/2$, then $d_G(x_j, i_j \pm 1) = 2$.

Assertion 3. For any j with $1 \leq j \leq m$, we have

(1) $d(i_j - 1) \leq n + 1 - |N(i_{j+1} - 1) \cup N(x_{j+1})| - \varepsilon$, where $i_{j+1} = i_2$ when $j = m$ and $i_1 = 1$; and $i_{j+1} = i_1$ when $j = m$ and $i_1 > 1$;

(2) $d(i_j + 1) \leq n + 1 - |N(i_{j-1} + 1) \cup N(x_{j-1})| - \varepsilon$, where $i_{j-1} = i_{m-1}$ when $j = 1$ and $i_m = t$ and $i_{j-1} = i_m$ when $j = 1$ and $i_m < t$, where $\varepsilon = 1$ for $2 \leq j \leq m-1$ and $(i_j - 1)(i_j + 1) \notin E(H)$, and $\varepsilon = 0$ for other cases,

Assertion 4. For any $u, v \in N^+$ or $u, v \in N^-$, $d(u) + d(v) \leq n$.

Assertion 5. For any $u \in N^- \cup N^+$, $d(u) \leq n/2$.

Assertion 6. For any j with $1 \leq j \leq m$, $|N(i_j \pm 1) \cup N(x_j)| \geq 2n/3$, if $i_j \pm 1 \in P$.

Assertion 7. For any j with $2 \leq j \leq m-1$, $(i_j - 1)(i_j + 1) \in E(H)$.

Thus using assertions 1-7, we can obtain a contradiction by a similar argument of Theorem 1.

Consider 3-connected graph $G = 3K_2 \vee K_3$, which is not hamilton-connected. It is easy to check that for all distinct vertices u and v of G with $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} < (n+1)/2$ and $|N(u) \cup N(v)| = 2n/3 - 1$. So, the conditions of Theorem 2 is best possible in a sense.

Corollary 2.1^[3] Let G be a 3-connected graph of order $n (\geq 3)$ and let u and v be distinct vertices of G . If $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq (n+1)/2$, then G is hamilton-connected.

Corollary 2.2^[4] Let G be a 3-connected graph of order $n (\geq 3)$, and let u and v be distinct vertices of G . If $d(u, v) = 2$ implies that $|N(u) \cup N(v)| \geq 2n/3$, then G is hamilton-connected.

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