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Further Results on Hilbert's Tenth Problem

Zhi-Wei Sun

Nanjing University
Nanjing 210093, P. R. China
zwsun@nju.edu.cn
<http://math.nju.edu.cn/~zwsun>

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Abstract

Hilbert's Tenth Problem (HTP) asked for an effective algorithm to test whether an arbitrary polynomial equation

$$P(x_1, \dots, x_n) = 0$$

(with integer coefficients) has solutions over the ring \mathbb{Z} of the integers. This was finally solved by Matiyasevich in 1970 negatively. In this talk we introduce the speaker's further results on HTP. In particular, we present a sketch of the proof of the speaker's main result that there is no effective algorithm to determine whether an arbitrary polynomial equation $P(x_1, \dots, x_{11}) = 0$ (with integer coefficients) in 11 unknowns has integral solutions or not.

Part I. Hilbert's Tenth Problem and its Solution

Hilbert's Tenth Problem

In 1900, at the Paris conference of ICM, D. Hilbert presented 23 famous mathematical problems. He formulated his tenth problem as follows:

Given a Diophantine equation with any number of unknown quantities and with rational integral numerical coefficients: To devise a process according to which it can be determined in a finite number of operations whether the equation is solvable in rational integers.

In modern language, Hilbert's Tenth Problem (HTP) asked for an effective algorithm to test whether an arbitrary polynomial equation

$$P(z_1, \dots, z_n) = 0$$

(with integer coefficients) has solutions over the ring \mathbb{Z} of the integers.

However, at that time the exact meaning of algorithm was not known.

The theory of computability

The theory of computability was born in the 1930's. A problem or a set is *decidable*, if and only if its characteristic function is Turing computable (or recursive).

An *r.e. (recursively enumerable) set* is the empty set \emptyset or the range of a recursive function, it is also the domain of a partial recursive function. An r.e. set is also said to be *semi-computable*.

A set $A \subseteq \mathbb{N} = \{0, 1, 2, \dots\}$ is recursive if and only if both A and $\mathbb{N} \setminus A$ are r.e. sets.

It is well-known that there are nonrecursive r.e. subsets of $\mathbb{N} = \{0, 1, 2, \dots\}$.

Diophantine relations and Diophantine sets

A relation $R(a_1, \dots, a_m)$ with $a_1, \dots, a_m \in \mathbb{N}$ is said to be *Diophantine* if there is a polynomial $P(t_1, \dots, t_m, x_1, \dots, x_n)$ with integer coefficients such that

$$R(a_1, \dots, a_m) \iff \exists x_1 \geq 0 \dots \exists x_n \geq 0 [P(a_1, \dots, a_m, x_1, \dots, x_n) = 0].$$

(Throughout this paper, variables always range over \mathbb{Z} .)

A set $A \subseteq \mathbb{N}$ is Diophantine if and only if the predicate $a \in A$ is Diophantine.

It is easy to see that any Diophantine set is an r.e. set.

Two key steps to solve HTP

In 1961 M. Davis, H. Putnam and J. Robinson [Ann. of Math.] successfully showed that any r.e. set is exponential Diophantine, that is, any r.e. set A has the exponential Diophantine representation

$$a \in A \iff \exists x_1 \geq 0 \dots \exists x_n \geq 0 [P(a, x_1, \dots, x_n, 2^{x_1}, \dots, 2^{x_n}) = 0],$$

where P is a polynomial with integer coefficients.

Recall that the Fibonacci sequence $(F_n)_{n \geq 0}$ defined by

$$F_0 = 0, F_1 = 1, \text{ and } F_{n+1} = F_n + F_{n-1} \quad (n = 1, 2, 3, \dots)$$

increases exponentially. In 1970 Yu. Matiyasevich took the last step to show ingeniously that the relation $y = F_{2^x}$ (with $x, y \in \mathbb{N}$) is Diophantine! It follows that the exponential relation $a = b^c$ (with $a, b, c \in \mathbb{N}$, $b > 1$ and $c > 0$) is Diophantine, i.e. there exists a polynomial $P(a, b, c, x_1, \dots, x_n)$ with integer coefficients such that

$$a = b^c \iff \exists x_1 \geq 0 \dots \exists x_n \geq 0 [P(a, b, c, x_1, \dots, x_n) = 0].$$

Matijasevich's theorem

Matijasevich's surprising result, together with the important work of Davis, Putnam and Robinson in 1961, leads to the following great result.

Matijasevich's Theorem (1970). Any r.e. set $A \subseteq \mathbb{N}$ is Diophantine.

As some r.e. sets are not recursive, it follows that there is *no* algorithm to decide whether an *arbitrary* polynomial equation

$$P(x_1, \dots, x_n) = 0$$

(with integer coefficients) has solutions $x_1, \dots, x_n \in \mathbb{N}$.

The negative solution to HTP

J. Robinson's Simple Observation:

$$\begin{aligned} & \exists z_1 \dots \exists z_n [P(z_1, \dots, z_n) = 0] \\ \iff & \exists x_1 \geq 0 \dots x_n \geq 0 \left[\prod_{\varepsilon_1, \dots, \varepsilon_n \in \{\pm 1\}} P(\varepsilon_1 x_1, \dots, \varepsilon_n x_n) = 0 \right]. \end{aligned}$$

On the other hand, by Lagrange's four-square theorem (each $m \in \mathbb{N}$ can be written as the sum of four squares), we have

$$\begin{aligned} & \exists x_1 \geq 0 \dots \exists x_n \geq 0 [P(x_1, \dots, x_n) = 0] \\ \iff & \exists u_1 \exists v_1 \exists y_1 \exists z_1 \dots \exists u_n \exists v_n \exists y_n \exists z_n \\ & [P(u_1^2 + v_1^2 + y_1^2 + z_1^2, \dots, u_n^2 + v_n^2 + y_n^2 + z_n^2) = 0] \end{aligned}$$

Therefore, the negative solution of HTP (over \mathbb{Z}) is equivalent to the negative solution of HTP (over \mathbb{N}).

Thus Matiyasevich solved HTP negatively!

Part II. Reduction of Natural Number Unknowns

Small ν with \exists^ν over \mathbb{N} undecidable

For a set $S \subseteq \mathbb{Z}$ we let \exists^n over S denote the set of formulas

$$\exists x_1 \in S \dots \exists x_n \in S [P(x_1, \dots, x_n) = 0]$$

with $P(x_1, \dots, x_n) \in \mathbb{Z}[x_1, \dots, x_n]$.

Any nonrecursive r.e. set A has a Diophantine representation:

$$a \in A \iff \exists x_1 \geq 0 \dots \exists x_n \geq 0 [P(x_1, \dots, x_n) = 0].$$

It is interesting to find the least $\nu \in \mathbb{Z}^+ = \{1, 2, 3, \dots\}$ such that \exists^ν over \mathbb{N} is undecidable.

$\nu < 200$ (Matiyasevich, Summer of 1970)

$\nu \leq 35$ (J. Robinson, 1970)

$\nu \leq 24$ (Matiyasevich and Robinson, 1970)

$\nu \leq 14$ (Matiyasevich and Robinson, 1970)

$\nu \leq 13$ (Matiyasevich and Robinson, 1973 [Acta Arith. 27(1975)])

$\nu \leq 9$ (Matiyasevich, 1975; details in Jones [J. Symbolic Logic, 1982])

Matiyasevich-Robinson's Relation-Combining Theorem

Matiyasevich-Robinson's Relation-Combining Theorem [Acta Arith. 27(1975)] Let A_1, \dots, A_k and R, S, T be integers with $S \neq 0$. Then

$$A_1 \in \square \wedge \dots \wedge A_k \in \square \wedge S \mid T \wedge R > 0 \\ \iff \exists n \geq 0 [M_k(A_1, \dots, A_k, S, T, R, n) = 0],$$

where

$$M_k(x_1, \dots, x_k, w, x, y, z) \\ = \prod_{\varepsilon_1, \dots, \varepsilon_k \in \{\pm 1\}} \left(x^2 + w^2 z - w^2(2y - 1) \left(x^2 + X^k + \sum_{j=1}^k \varepsilon_j \sqrt{x_j} X^{j-1} \right) \right) \\ = (w^2(1 - 2y))^{2k} J_k \left(x_1, \dots, x_k, x^2 + X^k + \frac{x^2 + w^2 z}{w^2(1 - 2y)} \right)$$

with $X = 1 + \sum_{j=1}^k x_j^2$, and $J_k(x_1, \dots, x_k, x)$ being

$$\prod_{\varepsilon_1, \dots, \varepsilon_k \in \{\pm 1\}} \left(x + \varepsilon_1 \sqrt{x_1} + \varepsilon_2 \sqrt{x_2} X + \dots + \varepsilon_k \sqrt{x_k} X^{k-1} \right).$$

My observation for later use

Matiyasevich-Robinson Relation-Combining Theorem is an important tool to reduce the number of unknowns.

Let $A_1, \dots, A_k \in \square$, and $R, S, T \in \mathbb{Z}$ with $R > 0$, $S \neq 0$ and $S \mid T$. We can easily see that

$$M_k(A_1, \dots, A_k, S, T, R, m) = 0,$$

where

$$\begin{aligned} m &= (2R - 1)(T^2 + X^k + \sqrt{A_1}X^0 + \dots + \sqrt{A_k}X^{k-1}) - \frac{T^2}{S^2} \\ &\geq X \geq \max\{A_1, \dots, A_k\} \end{aligned}$$

with $X = 1 + \sum_{j=1}^k A_j^2$.

A Useful Observation (Sun, 2017): If A_1, \dots, A_k are even and S is odd, then X is odd and

$$m \equiv (2R - 1)(T^2 + 1) - T^2 \equiv 1 \pmod{2}.$$

Coding idea of Matiyasevich and Robinson (1975)

Let $b \in \mathbb{N}$, $\delta \in \mathbb{Z}^+$, and

$$P(z_0, \dots, z_\nu) = \sum_{\substack{i_0, \dots, i_\nu \in \mathbb{N} \\ i_0 + \dots + i_\nu \leq \delta}} a_{i_0, \dots, i_\nu} z_0^{i_0} \cdots z_\nu^{i_\nu}.$$

$$B = 2\delta!(1 + b^\delta) \left(1 + \sum_{i_0 + \dots + i_\nu \leq \delta} a_{i_0, \dots, i_\nu}^2 \right) + 1,$$

$$D(x) = x^{(\delta+1)^{\nu+2}} + \sum_{i_0 + \dots + i_\nu \leq \delta} c_{i_0, \dots, i_\nu} a_{i_0, \dots, i_\nu} x^{(\delta+1)^{\nu+1} - \sum_{s=0}^{\nu} i_s (\delta+1)^s}$$

with $c_{i_0, \dots, i_\nu} = i_0! \dots i_\nu! (\delta - i_0 - \dots - i_\nu)!$. Then

$$P(z_0, \dots, z_\nu) = 0 \text{ for some } z_0, \dots, z_\nu \in [0, b]$$

\iff there is a number c of the form $1 + \sum_{i=0}^{\nu} c_i B^{(\delta+1)^i}$ with $c_i \in [0, b]$

such that the coefficient of $x^{(\delta+1)^{\nu+1}}$ in $(1 + \sum_{i=0}^{\nu} c_i x^{(\delta+1)^i})^\delta D(x)$

is zero.

Matiyasevich's idea to use binary representations

For $a, b \in \mathbb{N}$ written in base p with p prime, let $\tau_p(a, b)$ denote the number of carries occurring in the addition of a and b . Kummer noted that $\tau_p(a, b) = \text{ord}_p\left(\binom{a+b}{a}\right)$.

Let $b, B \in 2 \uparrow = \{2^n : n \in \mathbb{N}\}$ with $b \leq B$. Let $\delta, \nu \in \mathbb{Z}^+$. For $c = \sum_{j=0}^{(\delta+1)^\nu} c_j B^j$ with $c_j \in [0, B)$, and $M = \sum_{j=0}^{(\delta+1)^\nu} m_j B^j$ with

$$m_j = \begin{cases} B - b & \text{if } j = (\delta + 1)^s \text{ for some } s = 1, \dots, \nu, \\ B - 1 & \text{otherwise,} \end{cases}$$

$$\begin{aligned} \tau_2(c, M) = 0 &\iff \tau_2(c_j, m_j) = 0 \text{ for all } j = 0, \dots, (\delta + 1)^\nu \\ &\iff c = \sum_{i=1}^{\nu} z_i B^{(\delta+1)^i} \text{ for some } z_1, \dots, z_k \in [0, b) \end{aligned}$$

If $N \in 2 \uparrow$ and $S, T \in [0, N)$, then

$$\tau_2(S, T) = 0 \iff N^2 \mid \binom{2R}{R}$$

where $R = (N - 1)((S + T + 1)N + T + 1)$.

The 9 Unknowns Theorem

The above ideas, together with some other works in the 1975 paper of Matiyasevich and Robinson, led Matiyasevich obtain the following celebrated theorem.

Matiyasevich's 9 Unknowns Theorem: \exists^9 over \mathbb{N} is undecidable!

The detailed proof of this theorem appeared in Jones [J. Symbolic Logic, 1982].

Up to now, no one has shown that \exists^ν over \mathbb{N} is undecidable for some $\nu < 9$, although A.Baker, Matiyasevich and Robinson all believed that \exists^3 over \mathbb{N} might be undecidable.

Part III. Find small ν with \exists^ν over \mathbb{Z} undecidable

\exists over \mathbb{Z} is decidable

Matiyasevich and Robinson [Acta Arith. 27(1975)]: If a_0, a_1, \dots, a_n and z are integers with $a_0 z \neq 0$ and $\sum_{i=0}^n a_i z^{n-i} = 0$, then

$$|z|^n \leq |a_0 z^n| \leq \sum_{i=1}^n |a_i| |z|^{n-i} \leq \sum_{i=1}^n |a_i| |z|^{n-1}$$

and hence

$$|z| \leq \sum_{i=1}^n |a_i|.$$

Thus \exists over \mathbb{N} and \exists over \mathbb{Z} are decidable (in polynomial time).

It is not known whether \exists^2 over \mathbb{Z} is decidable. But A. Baker proved in 1968 that if $P(x, y) \in \mathbb{Z}[x, y]$ is homogenous, irreducible and of degree at least three then for any $m \in \mathbb{Z}$ there is an effective algorithm to determine whether $P(x, y) = m$ for some $x, y \in \mathbb{Z}$.

Relative results

For any $m \in \mathbb{Z}$, by Lagrange's four-square theorem

$$m \geq 0 \iff \exists z_1 \exists z_2 \exists z_3 \exists z_4 [m = z_1^2 + z_2^2 + z_3^2 + z_4^2].$$

Thus

$$\exists^n \text{ over } \mathbb{N} \text{ is undecidable} \Rightarrow \exists^{4n} \text{ over } \mathbb{Z} \text{ is undecidable.}$$

By the Gauss-Legendre theorem on sums of three squares,

$$\mathbb{N} \setminus \{x^2 + y^2 + z^2 : x, y, z \in \mathbb{Z}\} = \{4^k(8l + 7) : k, l \in \mathbb{N}\}.$$

If $n \in \mathbb{N}$, then $4n + 1 = (2x)^2 + (2y)^2 + (2z + 1)^2$ for some $x, y, z \in \mathbb{Z}$, and hence $n = x^2 + y^2 + z^2 + z$. Thus, for any $m \in \mathbb{Z}$,

$$m \geq 0 \iff \exists x \exists y \exists z [m = x^2 + y^2 + z^2 + z].$$

It follows that

$$\exists^n \text{ over } \mathbb{N} \text{ is undecidable} \Rightarrow \exists^{3n} \text{ over } \mathbb{Z} \text{ is undecidable.}$$

Thus \exists^{27} over \mathbb{Z} is undecidable by the 9 unknowns theorem, as pointed out by S.P. Tung in [Japan J. Math., 11(1985)].

A new relation-combining theorem

Tung (1985) asked whether \exists^ν over \mathbb{Z} is undecidable for some $\nu < 27$.

New Relation-Combining Theorem (Z.-W. Sun [Z. Math. Logik Grundlag. Math. 38(1992)]): Let $A_1, \dots, A_k, B, C_1, \dots, C_n, D, E$ be integers with $D \neq 0$. Then

$$A_1, \dots, A_k \in \square \wedge B \neq 0 \wedge C_1, \dots, C_n \geq 0 \wedge D \mid E \\ \iff \exists z_1 \dots \exists z_{n+2} [P(A_1, \dots, A_k, B, C_1, \dots, C_n, D, E, z_1, \dots, z_{n+2}) = 0],$$

where P is a suitable polynomial with integer coefficients.

This implies that

$$\exists^n \text{ over } \mathbb{N} \text{ is undecidable} \Rightarrow \exists^{2n+2} \text{ over } \mathbb{Z} \text{ is undecidable.}$$

So \exists^{20} over \mathbb{Z} is undecidable by the 9 unknowns theorem.

\exists^{11} over \mathbb{Z} is undecidable

In 1992, I announced that \exists^{11} **over \mathbb{Z} is undecidable.**

To achieve this goal, unlike others I did not simply use the relative result, instead I adapt the deep proof of the 9 unknowns theorem and made suitable variants so that we can use integer variables instead of natural number variables.

My starting point is the use of Lucas sequences with integer indices instead of the usual natural number indices. I published this initial step in Sci. China Ser. A 35(1992).

The whole proof of the undecidability of \exists^{11} over \mathbb{Z} is very sophisticated. It appeared in my PhD thesis in 1992. During 1992-2016, despite that many mathematicians wanted to see my detailed proof, I did not write an English version of that, since I was frequently busy with my new discoveries.

After 25 years have passed, I finally spent time to write an English paper which contains the undecidability of \exists^{11} over \mathbb{Z} as well as my new discoveries related to HTP. The preprint is now publicly available from <http://arxiv.org/abs/1704.03504>

Lucas sequences

Let A and B be integers. The usual Lucas sequence $u_n = u_n(A, B)$ ($n = 0, 1, 2, \dots$) and its companion $v_n = v_n(A, B)$ ($n = 0, 1, 2, \dots$) are defined as follows:

$$u_0 = 0, u_1 = 1, \text{ and } u_{n+1} = Au_n - Bu_{n-1} \quad (n = 1, 2, 3, \dots);$$

and

$$v_0 = 2, v_1 = 2, \text{ and } v_{n+1} = Av_n - Bv_{n-1} \quad (n = 1, 2, 3, \dots).$$

Note that

$$u_n(2, 1) = n, u_n(1, -1) = F_n, \text{ and } u_n(3, 1) = F_{2n}.$$

Let

$$\alpha = \frac{A + \sqrt{\Delta}}{2} \text{ and } \beta = \frac{A - \sqrt{\Delta}}{2}$$

be the two roots of the quadratic equation $x^2 - Ax + B = 0$ where $\Delta = A^2 - 4B$. It is well known that for any $n \in \mathbb{N}$ we have

$$(\alpha - \beta)u_n = \alpha^n - \beta^n, v_n = \alpha^n + \beta^n \text{ and } v_n^2 - \Delta u_n^2 = 4B^n.$$

Pell's equation

Let $d \in \mathbb{Z}^+ \setminus \square$. It is well-known that the Pell equation

$$y^2 - dx^2 = 1$$

has infinitely many integral solutions. (Note that $x = 0$ and $y = \pm 1$ are trivial solutions.) Moreover,

$$\{y + \sqrt{d}x : x, y \in \mathbb{Z} \text{ and } y^2 - dx^2 = 1\}$$

is a multiplicative cyclic group.

For any integer $A \geq 2$, the solutions of the Pell equation

$$y^2 - (A^2 - 1)x^2 = 1 \quad (x, y \in \mathbb{N})$$

are given by $x = u_n(2A, 1)$ and $y = v_n(2A, 1)$ with $n \in \mathbb{N}$. J. Robinson and his followers wrote $u_n(2A, 1)$ and $v_n(2A, 1)$ as $\psi_n(A)$ and $\chi_n(A)$ respectively.

To unify Matiyasevich's use of $F_{2n} = u_n(3, 1)$ and Robinson's use of $\psi_n(A) = u_n(2A, 1)$, we deal with Lucas sequences $(u_n(A, 1))_{n \geq 0}$.

On $u_n(A, 1)$ with $n \in \mathbb{Z}$

We extend the sequences $u_n = u_n(A, 1)$ and $v_n = v_n(A, 1)$ to integer indices by letting

$$u_0 = 0, \quad u_1 = 1, \quad \text{and} \quad u_{n-1} + u_{n+1} = Au_n \quad \text{for all } n \in \mathbb{Z},$$

and

$$v_0 = 2, \quad v_1 = A, \quad \text{and} \quad v_{n-1} + v_{n+1} = Av_n \quad \text{for all } n \in \mathbb{Z}.$$

It is easy to see that

$$u_{-n}(A, 1) = -u_n(A, 1) = (-1)^n u_n(-A, 1)$$

and $v_{-n}(A, 1) = v_n(A, 1) = (-1)^n v_n(-A, 1)$ for all $n \in \mathbb{Z}$.

Lemma. Let $A, X \in \mathbb{Z}$. Then

$$(A^2 - 4)X^2 + 4 \in \square \iff X = u_m(A, 1) \quad \text{for some } m \in \mathbb{Z}.$$

Remark. For $n \in \mathbb{N}$ and $A \geq 2$, it is easy to show that

$$(A - 1)^n \leq u_{n+1}(A, 1) \leq A^n.$$

Diophantine representation of $C = u_B(A, 1)$ with unknowns arbitrarily large

Matiyasevič and Robinson (1975) showed that for $A > 1$ and $B, C > 0$ there is a Diophantine representation of $C = u_B(2A, 1)$ only involving three natural number variables.

Lemma (Sun [Sci. China Ser. A 35(1992)]). Let $A, B, C \in \mathbb{Z}$ with $A > 1$ and $B \geq 0$. Then

$$C = u_B(A, 1) \iff C \geq B \wedge \exists x > 0 \exists y > 0 (DFI \in \square) \\ \iff \exists x, y, z \geq 0 [DFI(C - B + 1)^2 = (z - DFI(C - B + 1))^2],$$

where

$$D = (A^2 - 4)C^2 + 4, \quad E = C^2 D x, \quad F = 4(A^2 - 4)E^2 + 1, \\ G = 1 + CDF - 2(A + 2)(A - 2)^2 E^2, \quad H = C + BF + (2y - 1)CF, \\ I = (G^2 - 1)H^2 + 1.$$

Moreover, if $C = u_B(A, 1)$ with $B > 0$, then for any $Z \in \mathbb{Z}^+$ there are integers $x \geq Z$ and $y \geq Z$ with $DFI \in \square$.

Diophantine representation of $C = u_B(A, 1)$ with integer unknowns

Clearly $C \geq B \iff \exists x \geq 0 (C = B + x)$. However, if we use integer variables, we need three variables:

$$C \geq B \iff \exists x \exists y \exists z [C = B + x^2 + y^2 + z^2 + z].$$

Thus, to save the number of integer variables involved, we should try to avoid inequalities.

Note that

$$u_B(A, 1) \equiv u_B(2, 1) = B \pmod{A - 2}.$$

Lemma (Sun [Sci. China Ser. A 35(1992)]). Let $A, B, C \in \mathbb{Z}$ with $1 < |B| < |A|/2 - 1$. Then

$$C = u_B(A, 1) \iff (A - 2 \mid C - B) \wedge \exists x \neq 0 \exists y (DFI \in \square),$$

where D, F, I are defined as before.

Diophantine representation of $W = V^B$ with integer unknowns

J. Robinson showed that $W = V^B$ (with $V > 1$ and $B, W > 0$) if and only if there is an integer $A > \max\{V^{3B}, W^B\}$ such that

$$(V^2 - 1)W u_B(2A, 1) \equiv V(W^2 - 1) \pmod{2AV - V^2 - 1}.$$

Lemma (Sun [Sci. China Ser. A 35(1992)]). Let B, V, W be integers with $B > 0$ and $|V| > 1$. Then $W = V^B$ if and only if there are $A, C \in \mathbb{Z}$ for which $|A| \geq \max\{V^{4B}, W^4\}$, $C = u_B(A, 1)$ and

$$(V^2 - 1)WC \equiv V(W^2 - 1) \pmod{AV - V^2 - 1}.$$

Remark. A, V and W in this lemma are not necessarily positive, they might be negative. We have also shown that for $B, V, W \in \mathbb{Z}$ with $B > 0$ and $|V| > 1$, the equality $W = V^B$ holds if and only if there are integers A and C for which $|A| \geq \max\{V^{2B}, W^2\}$, $C = u_{2B+1}(A, 1)$ and

$$(V - 1)WC \equiv VW^2 - 1 \pmod{(A^2 - 2)V - V^2 - 1}.$$

The first auxiliary theorem

Theorem 1 (Sun, arXiv:1704.03504). Let $\mathcal{A} \subseteq \mathbb{N}$ be a Diophantine set, and let p be a prime. Then, for each $a \in \mathbb{N}$, we have

$$a \in \mathcal{A} \Rightarrow \forall Z > 0 \exists f \geq Z \exists g \in [b, C) \left(b \in \square \wedge b \in p \uparrow \wedge Y \mid \binom{pX}{X} \right)$$

and

$$\exists f \neq 0 \exists g \in [0, 2C) \left(b \in \square \wedge b \in p \uparrow \wedge Y \mid \binom{pX}{X} \right) \Rightarrow a \in \mathcal{A},$$

where

$$b := 1 + (p^2 - 1)(ap + 1)f,$$

$C = p^{\alpha_1 p} b^{\alpha_2}$ for some $\alpha_1, \alpha_2 \in \mathbb{Z}^+$ only depending on \mathcal{A} , and X and Y are suitable polynomials in $\mathbb{Z}[a, f, g]$ such that if $a \in \mathbb{N}$, $f \in \mathbb{Z} \setminus \{0\}$, $b \in \square$ and $0 \leq g < 2C$ then

$$p + 1 \mid X, \quad X \geq 3b \quad \text{and} \quad Y \geq \max\{b, p^{4p}\}.$$

Remark. Clearly, $b \in \square \wedge f \neq 0 \Rightarrow f > 0 \wedge b > a \wedge p^2 - 1 \mid b - 1$.

The second auxiliary theorem

Theorem 2 (Sun, arXiv:1704.03504). Let p be a prime, and let $b \in p \uparrow$ and $g \in \mathbb{Z}^+$. Let P, Q, X, Y be integers with $P > Q > 0$ and $X, Y \geq b$. Suppose that $Y \mid \binom{PX}{QX}$. Then there are integers $h, k, l, w, x, y \geq b$ for which

$$DFI \in \square, (U^{2P}V^2 - 4)K^2 + 4 \in \square, pA - p^2 - 1 \mid (p^2 - 1)WC - p(W^2 - 1),$$
$$bw = p^B \text{ and } 16g^2(C - KL)^2 < K^2,$$

where

$$L := lY, U := PLX, V := 4gwY,$$

$$W := bw, K := QX + 1 + k(U^P V - 2),$$

$$A := U^Q(V + 1), B := PX + 1, C := B + (A - 2)h,$$

and D, F, I are as before.

Remark. We actually take $C = u_B(A, 1)$, $K = u_{QX+1}(U^P V, 1)$,
 $L = \lfloor (V + 1)^{PX} / V^{QX} \rfloor \equiv \binom{PX}{QX} \pmod{V}$.

The third auxiliary theorem

Theorem 3 (Sun, arXiv:1704.03504). Let p be a prime, and let $b \in \mathbb{N}$ and $g \in \mathbb{Z}^+$. Let P, Q, X, Y be integers with

$$P > Q > 0, \quad X \geq 3b, \quad \text{and} \quad Y \geq \max\{b, p^{4P}\}.$$

Suppose that there are integers h, k, l, w, x, y with $lx \neq 0$ such that

$$DFI \in \square, \quad (U^{2P}V^2 - 4)K^2 + 4 \in \square, \quad pA - p^2 - 1 \mid (p^2 - 1)WC - p(W^2 - 1),$$

and

$$4(C - KL)^2 < K^2,$$

where we adopt previous notations. Then

$$b \in p \uparrow \quad \text{and} \quad Y \mid \begin{pmatrix} PX \\ QX \end{pmatrix}.$$

Remark. This theorem involving integer variables plays a central role in our proof of the undecidability of \exists^{11} over \mathbb{Z} .

Outline of the proof (I)

First show that $bw = W \neq 0$ and $1 < B < |A|/2 - 1$. As $x \neq 0$ and $DFI \in \square$, we obtain $C = u_B(A, 1)$.

Since $((U^P V)^2 - 4)K^2 + 4 \in \square$, we have $K = u_R(U^P V, 1)$ for some $R \in \mathbb{Z}$. Note that

$$QX + 1 \equiv K = u_R(U^P V, 1) \equiv u_R(2, 1) = R \pmod{U^P V - 2}.$$

If $R \neq QX + 1$, then we get

$$\left| \frac{C}{K} \right| = \left| \frac{u_{|R|+1}(|A|, 1)}{u_{|R|}(|U^P V|, 1)} \right| \leq \left(\frac{|U^P|}{|U^P V| - 1} \right)^{QX} < \left(\frac{1}{2} \right)^{QX} \leq \frac{1}{2}$$

and hence

$$|L| \leq \left| L - \frac{C}{K} \right| + \left| \frac{C}{K} \right| < \frac{1}{2} + \frac{1}{2} \leq 1$$

which contradicts $L = lY \neq 0$.

Outline of the proof (II)

Now, $R = QX + 1$ and $K = u_{QX+1}(U^P V, 1)$. Set

$$\rho = \frac{(V+1)^{PX}}{V^{QX}}.$$

Then we can show $|C/K| \geq |\rho|/2$. Since

$$\frac{|V+1|^{Q+1}}{|V|^Q} \geq \frac{|V|-1}{2},$$

we have $|\rho| \geq ((|V|-1)/2)^X \geq (2Q)^X \geq 2^X \geq 2$. Thus

$$|L| > \left| \frac{C}{K} \right| - \frac{1}{2} \geq \frac{|\rho|}{2} - \frac{1}{2} \geq \frac{|\rho|}{4} \geq \frac{1}{4} \left(\frac{|V|-1}{2} \right)^X,$$

$$|A| \geq \left(\frac{|V|-1}{2} \right)^{X+1} \geq \max\{p^{4B}, W^4\}.$$

Since $C = u_B(A, 1)$ and

$$(p^2 - 1)WC \equiv p(W^2 - 1) \pmod{pA - p^2 - 1},$$

we obtain $bw = W = p^B = p^{PX+1}$ and hence $b, w \in p \uparrow$.

Outline of the proof (III)

Now, $V = 4gwY \geq 4gwb = 4gW = 4p^{PX+1} \geq 8 \times 2^{PX}$, and hence

$$0 \leq \frac{1}{V} \sum_{i=0}^{QX-1} \frac{\binom{PX}{i}}{V^{QX-1-i}} < \frac{2^{PX}}{V} \leq \frac{1}{8}.$$

Note that

$$\rho = \frac{(V+1)^{PX}}{V^{QX}} = \frac{1}{V} \sum_{i=0}^{QX-1} \frac{\binom{PX}{i}}{V^{QX-1-i}} + \binom{PX}{QX} + \sum_{i=QX+1}^{PX} \binom{PX}{i} V^{i-QX},$$

and Y divides both L and V . So $\lfloor \rho \rfloor = L \Rightarrow Y \mid \binom{PX}{QX}$.

If $|C/K - \rho| < 1/4$, then

$$|\lfloor \rho \rfloor - L| \leq |\lfloor \rho \rfloor - \rho| + \left| \rho - \frac{C}{K} \right| + \left| \frac{C}{K} - L \right| < \frac{1}{8} + \frac{1}{4} + \frac{1}{2} < 1$$

and hence $\lfloor \rho \rfloor = L$ as desired.

Outline of the proof (IV)

Observe that

$$A^{PX} C = A^{PX} u_{PX+1}(A, 1) = |A|^{PX} u_{PX+1}(|A|, 1) > 0,$$

$$\begin{aligned}(U^P V)^{QX} K &= (U^P V)^{QX} u_{QX+1}(U^P V, 1) \\ &= |U^P V|^{QX} u_{QX+1}(|U^P V|, 1) > 0,\end{aligned}$$

and

$$A^{PX} (U^P V)^{QX} = U^{2PQX} (V + 1)^{PX} V^{QX} > 0.$$

So $CK > 0$ and hence

$$\left| \frac{C}{K} - \rho \right| = \left| \left| \frac{C}{K} \right| - |\rho| \right| \leq \rho \frac{2PX}{|U|V} = \frac{\rho}{|L|} \times \frac{2}{V} < \frac{8}{V} \leq \frac{1}{2^{PX}} \leq \frac{1}{4}$$

as desired.

Main Theorem

Theorem (Sun, arXiv:1704.03504). Let $\mathcal{A} \subseteq \mathbb{N}$ be an r.e. set.

(i) There is a polynomial $P_{\mathcal{A}}(z_0, z_1, \dots, z_9)$ with integer coefficients such that for any $a \in \mathbb{N}$ we have

$$\exists z_1 \dots \exists z_8 \exists z_9 \geq 0 [P_{\mathcal{A}}(a, z_1, \dots, z_9) = 0] \implies a \in \mathcal{A},$$

and

$$a \in \mathcal{A} \implies \forall Z > 0 \exists z_1 \geq Z \dots \exists z_9 \geq Z [P_{\mathcal{A}}(a, z_1, \dots, z_9) = 0].$$

(ii) There is a polynomial $Q_{\mathcal{A}}(z_0, z_1, \dots, z_{10})$ with integer coefficients such that for any $a \in \mathbb{N}$ we have

$$a \in \mathcal{A} \iff \exists z_1 \dots \exists z_9 \exists z_{10} \neq 0 [Q_{\mathcal{A}}(a, z_1, \dots, z_{10}) = 0].$$

Proof of the first part (I)

By Matiyasevich's theorem, \mathcal{A} is a Diophantine set. Let p be a prime, and let b, \mathcal{C} and $X, Y \in \mathbb{Z}[a, f, g]$ be as in the first auxiliary theorem. Set $P = p$ and $Q = 1$, and adopt previous notations for capital Latin letters.

(i) Suppose that $a \in \mathcal{A}$. For any $Z \in \mathbb{Z}^+$ we may take $f \geq Z$ with $b \in \square$ and $b \in p \uparrow$, and $g \in [b, \mathcal{C})$ with Y dividing $\binom{PX}{QX} = \binom{pX}{X}$. Clearly,

$$0 < f \leq b \leq g < \mathcal{C} < 2\mathcal{C}.$$

By the second auxiliary theorem, there are integers $h, k, l, w, x, y \geq b$ such that

$$DFI \in \square, (U^{2P}V^2 - 4)K^2 + 4 \in \square, pA - p^2 - 1 \mid (p^2 - 1)WC - p(W^2 - 1),$$

(*)

$$bw = p^B \quad \text{and} \quad 16g^2(C - KL)^2 < K^2,$$

Thus

$$4(C - KL)^2 + \frac{g^2 K^2}{8\mathcal{C}^3} < \frac{K^2}{4g^2} + \frac{K^2}{8g} \leq \frac{K^2}{g}.$$

Proof of the first part (II)

Hence

$$O := f^2 l^2 x^2 (8C^3 g K^2 - g^2 (32(C - KL)^2 C^3 + g^2 K^2)) > 0.$$

Note that

$$g, h, k, l, w, x, y \geq b \geq f \geq Z.$$

By Matiyasevich-Robinson Relation-Combining Theorem we have

$$P_{\mathcal{A}}(a, f, g, h, k, l, w, x, y, m) = 0$$

for some integer $m \geq b \geq f \geq Z$, where

$$\begin{aligned} & P_{\mathcal{A}}(a, f, g, h, k, l, m, w, x, y) \\ &= M_3(b, DfI, (U^{2P} V^2 - 4)K^2 + 4, pA - p^2 - 1, \\ & \quad (p^2 - 1)WC - p(W^2 - 1), O, m). \end{aligned}$$

Note that $P_{\mathcal{A}}(z_0, z_1, \dots, z_9) \in \mathbb{Z}[z_0, z_1, \dots, z_9]$.

Proof of the first part (III)

Let $a \in \mathbb{N}$, and assume that there are integers $m \geq 0$ and f, g, h, k, l, w, x, y such that

$$P_{\mathcal{A}}(a, f, g, h, k, l, w, x, y, m) = 0.$$

By the Relation-Combining Theorem,

$$DFI \in \square, (U^{2P}V^2 - 4)K^2 + 4 \in \square, pA - p^2 - 1 \mid (p^2 - 1)WC - p(W^2 - 1),$$

(*)

also $b \in \square$ and

$$O := f^2 l^2 x^2 (8C^3 g K^2 - g^2 (32(C - KL)^2 C^3 + g^2 K^2)) > 0.$$

Clearly, $fglx \neq 0$. As $b \geq 0$ and $f \neq 0$, we have $b > 0$ and hence $C > 0$. Observe that

$$\frac{K^2}{g} > 4(C - KL)^2 + \frac{g^2 K^2}{8C^3} \geq \frac{g^2 K^2}{8C^3} \geq 0.$$

Thus $K \neq 0$ and $0 < g < 2C$. By the third auxiliary theorem, we have $b \in p \uparrow$ and $\begin{pmatrix} pX \\ X \end{pmatrix} = \begin{pmatrix} PX \\ QX \end{pmatrix} \equiv 0 \pmod{Y}$. Hence $a \in \mathcal{A}$ by the first auxiliary theorem.

Two Lemmas

Lemma 1. For any $A_1, \dots, A_k, S, T \in \mathbb{Z}$ with $S \neq 0$, we have

$$A_1 \in \square \wedge \dots \wedge A_k \in \square \wedge S \mid T \iff \exists z [H_k(A_1, \dots, A_k, S, T, z) = 0],$$

where

$$H_k(x_1, \dots, x_k, x, y, z) := x^{2k} J_k \left(x_1, \dots, x_k, z - \frac{y}{x} \right).$$

Remark. This is motivated by Matiyasevich-Robinson's Relation-Combining Theorem. Note that z is an integer variable.

Lemma 2 (Sun, arXiv:1704.03504). Let $m \in \mathbb{Z}$. Then

$$m \geq 0 \iff \exists x \neq 0 [(3m - 1)x^2 + 1 \in \square].$$

Remark. This is easy since if $m \in \mathbb{Z}^+$ then $3m - 1 \notin \square$ and hence the Pell equation

$$y^2 - (3m - 1)x^2 = 1$$

has infinitely many integral solutions.

Proof of the second part

Proof of the Second Part of the Main Theorem. A nonnegative integer a belongs to \mathcal{A} , if and only if there are integers f, g, h, k, l, w, x, y such that $b \in \square$, $O > 0$ and $(*)$ holds. Clearly,

$$O > 0 \iff O - 1 \geq 0 \iff \exists z \neq 0 [(3O - 4)z^2 + 1 \in \square].$$

In light of Lemma 1, we have

$$\begin{aligned} & b \in \square, (3O - 4)z^2 + 1 \in \square, \text{ and } (*) \text{ holds} \\ \iff & \exists m [Q_{\mathcal{A}}(a, f, g, h, k, l, m, w, x, y, z) = 0], \end{aligned}$$

where

$$\begin{aligned} & Q_{\mathcal{A}}(a, f, g, h, k, l, m, w, x, y, z) \\ = & H_4(b, (3O - 4)z^2 + 1, DFI, (U^{2P}V^2 - 4)K^2 + 4, \\ & pA - p^2 - 1, (p^2 - 1)WC - p(W^2 - 1), m). \end{aligned}$$

Note that $Q_{\mathcal{A}}(z_0, z_1, \dots, z_{10}) \in \mathbb{Z}[z_0, z_1, \dots, z_{10}]$.

Corollary 1

As some r.e. sets are not Diophantine, the Main Theorem has the following consequence.

Corollary 1. (i) There is no algorithm to determine for any $P(z_1, \dots, z_9) \in \mathbb{Z}[z_1, \dots, z_9]$ whether the equation

$$P(z_0, \dots, z_9) = 0$$

has integral solutions with $z_9 \geq 0$ (or $z_1 + \dots + z_9 \geq 0$).

(ii) There is no algorithm to determine for any $Q(z_1, \dots, z_{10}) \in \mathbb{Z}[z_1, \dots, z_9]$ whether the equation

$$Q(z_0, \dots, z_{10}) = 0$$

has integral solutions with $z_{10} \neq 0$ (or $z_1 + \dots + z_{10} \neq 0$).

Remark. Let $z'_9 = z_9 - z_1 - \dots - z_8$. Then

$$\begin{aligned} P(z_1, \dots, z_8, z'_9) = 0 \text{ with } z_1 + \dots + z_8 + z'_9 \geq 0 \\ \iff P(z_1, \dots, z_8, z_9) = 0 \text{ with } z_9 \geq 0. \end{aligned}$$

\exists^{11} over \mathbb{Z} is undecidable

Recall that

$$m \geq 0 \iff \exists x \exists y \exists z [m = x^2 + y^2 + z^2 + z].$$

So,

$$\begin{aligned} & \exists z_1 \dots \exists z_8 \exists z_9 \geq 0 [P(z_1, \dots, z_8, z_9) = 0] \\ \iff & \exists z_1 \dots \exists z_{11} [P(z_1, \dots, z_8, z_9^2 + z_{10}^2 + z_{11}^2 + z_{11}) = 0]. \end{aligned}$$

Similarly, in view of S.P. Tung's observation (1985)

$$m \neq 0 \iff \exists x \exists y [m = (2x + 1)(2y + 1)],$$

we have

$$\begin{aligned} & \exists z_1 \dots \exists z_9 \exists z_{10} \neq 0 [Q(z_1, \dots, z_9, z_{10}) = 0] \\ \iff & \exists z_1 \dots \exists z_{11} [Q(z_1, \dots, z_9, (2z_{10} + 1)(3z_{11} + 1)) = 0]. \end{aligned}$$

Therefore, both parts of the Main Theorem implies the undecidability of \exists^{11} over \mathbb{Z} .

Corollary 2

By taking negations of the formulas

$$\exists z_1 \dots \exists z_8 \exists z_9 \geq 0 [P(z_1, \dots, z_9) = 0]$$

and

$$\exists z_1 \dots \exists z_9 \exists z_{10} \neq 0 [Q(z_1, \dots, z_{10}) = 0]$$

in the Main Theorem, we get the following result.

Corollary 2 (Sun, arXiv:1704.03504) (i) $\forall^9 \exists^3$ over \mathbb{Z} is undecidable, i.e., there is no algorithm to test whether

$$\forall z_1 \dots \forall z_9 \exists x \exists y \exists z [P(z_1, \dots, z_9, x, y, z) = 0],$$

where P is an arbitrary polynomial of 12 variables with integer coefficients.

(ii) $\forall^{10} \exists^2$ over \mathbb{Z} is undecidable, i.e., there is no algorithm to test whether

$$\forall z_1 \dots \forall z_{10} \exists x \exists y [Q(z_1, \dots, z_{10}, x, y) = 0],$$

where Q is an arbitrary polynomial of 12 variables with integer coefficients.

Quantifier prefixes over Diophantine equations

In 1987 S.P. Tung proved for each $n \in \mathbb{Z}^+$ that $\forall^n \exists$ over \mathbb{Z} is co-NP-complete. He also showed that $\forall^{27} \exists^2$ over \mathbb{Z} is undecidable, and asked whether 27 here can be replaced by a smaller number. Corollary 2 of us tells that $\forall^{10} \exists^2$ over \mathbb{Z} and $\forall^9 \exists^3$ over \mathbb{Z} are undecidable.

In 1975 Matiyasevich and Robinson showed that $\exists^2 \forall \exists$ with \forall bounded is undecidable over \mathbb{N} . In 1981 Jones obtained the decidability of $\forall \exists$ over \mathbb{N} as well as some other undecidable results over \mathbb{N} .

In my PhD thesis in 1992, I also proved that

$$\forall \exists^6, \forall^2 \exists^4, \forall \exists \forall \exists^3, \forall \exists \forall^3 \exists^2, \forall^2 \exists \forall^2 \exists^2, \forall \exists^2 \forall^2 \exists^2, \\ \exists^2 \forall \exists^3, \exists^2 \forall^3 \exists^2, \exists \forall \exists \forall^2 \exists^2, \exists \forall \exists^4, \exists \forall^2 \exists^3, \exists \forall^5 \exists^2$$

over \mathbb{Z} are undecidable, and that

$$\exists^2 \forall \exists^3, \exists^2 \forall^2 \exists^2, \exists \forall \exists \forall \exists^2, \exists \forall \exists^4, \exists \forall^2 \exists^3, \exists \forall^4 \exists^2$$

with \forall bounded by polynomials are undecidable over \mathbb{Z} .

Part IV. Undecidable results related to polygonal numbers

Polygonal numbers

Recall that triangular numbers have the form $T_x = x(x+1)/2$ with $x \in \mathbb{Z}$. Note that $T_{-1-x} = T_x$.

Polygonal numbers are nonnegative integers constructed geometrically from the regular polygons. For $m = 3, 4, 5, \dots$, the m -gonal numbers are given by

$$p_m(n) = (m-2) \binom{n}{2} + n \quad (n = 0, 1, 2, \dots).$$

Clearly

$$p_3(n) = T_n, \quad p_4(n) = n^2, \quad p_5(n) = \frac{3n^2 - n}{2}, \quad p_6(n) = 2n^2 - n = T_{2n-1}.$$

The larger m is, the more sparse m -gonal numbers are.

Fermat claimed that for each $m = 3, 4, \dots$ any $n \in \mathbb{N}$ can be written as the sum of m polygonal numbers of order m . This was proved by Lagrange for $m = 4$, Gauss for $m = 3$, and Cauchy for $m \geq 5$.

Generalized pentagonal numbers and octagonal numbers

For $m = 5, 6, \dots$ those $p_m(x)$ with $x \in \mathbb{Z}$ are called *generalized polygonal numbers of order m* . We set

$$\text{Tri} = \{T_x : x \in \mathbb{Z}\}, \text{ Pen} = \left\{ p_5(x) = \frac{x(3x-1)}{2} : x \in \mathbb{Z} \right\}$$

and

$$\text{Octa} = \{p_8(x) = x(3x-2) : x \in \mathbb{Z}\}.$$

R. K. Guy [Amer. Math. Monthly 101(1994)]: Each $n \in \mathbb{N}$ is the sum of three elements of Pen.

Z.-W. Sun [J. Number Theory, 162(2016)]: Any $n \in \mathbb{N}$ is the sum of four elements of Octa. (This is quite similar to Lagrange's four-square theorem.)

Clearly,

$$x = \frac{x(x+1)}{2} - \frac{x(x-1)}{2} = T_x - T_{-x},$$

$$x = \frac{x(3x+1)}{2} - \frac{x(3x-1)}{2} = p_5(-x) - p_5(x).$$

A lemma on squares and generalized octagonal numbers

Lemma (Sun, arXiv:1704.03504). (i) Any integer can be written as $2^\delta(x^2 - y^2)$ with $\delta \in \{0, 1\}$ and $x, y \in \mathbb{Z}$. Also, each integer can be written as $2^\delta(p_8(x) - p_8(y))$ with $\delta \in \{0, 1\}$ and $x, y \in \mathbb{Z}$.
(ii) Any positive odd integer can be written as $x^2 + y^2 + 2z^2$ with $x, y, z \in \mathbb{Z}$. Also, each positive odd integer can be written as $p_8(x) + p_8(y) + 2p_8(z)$ with $x, y, z \in \mathbb{Z}$.

The first assertion in part (ii) is known.

Let $n \in \mathbb{Z}^+$. By Lemma 4.3(ii) of Sun [J. Number Theory, 162(2016)], $6n + 1 = x^2 + y^2 + 2z^2$ for some $x, y, z \in \mathbb{Z}$ with $3 \nmid xyz$. (This is a nontrivial result!) Without loss of generality we may assume that $x = 3u - 1$, $y = 3v - 1$ and $z = 3w - 1$ for some $u, v, w \in \mathbb{Z}$. Thus

$$\begin{aligned}6n + 1 &= (3u - 1)^2 + (3v - 1)^2 + 2(3w - 1)^2 \\ &= (3p_8(u) + 1) + (3p_8(v) + 1) + 2(3p_8(w) + 1)\end{aligned}$$

and hence $2n - 1 = p_8(u) + p_8(v) + 2p_8(w)$.

Undecidable results related to Tri, \square , Pen and Octa

Theorem (Z. W. Sun, arXiv:1704.03504). Let \mathcal{A} be any r.e. subset of \mathbb{N} . Then there is a polynomial $P_4(z_0, z_1, \dots, z_{17})$ with integer coefficients such that for any $a \in \mathbb{N}$ we have

$$a \in \mathcal{A} \iff \exists z_1 \in \square \dots \exists z_{17} \in \square [P_4(a, z_1, \dots, z_{17}) = 0].$$

Also, there are polynomials

$$P_3(z_0, z_1, \dots, z_{18}), P_5(z_0, z_1, \dots, z_{18}), P_8(z_0, z_1, \dots, z_{18})$$

with integer coefficients such that for any $a \in \mathbb{N}$ we have

$$\begin{aligned} a \in \mathcal{A} &\iff \exists z_1 \in \text{Tri} \dots \exists z_{18} \in \text{Tri} [P_3(a, z_1, \dots, z_{18}) = 0] \\ &\iff \exists z_1 \in \text{Pen} \dots \exists z_{18} \in \text{Pen} [P_5(a, z_1, \dots, z_{18}) = 0] \\ &\iff \exists z_1 \in \text{Octa} \dots \exists z_{18} \in \text{Octa} [P_8(a, z_1, \dots, z_{18}) = 0], \end{aligned}$$

Corollary. \exists^{17} over \square , \exists^{18} over Tri, \exists^{18} over Pen, and \exists^{18} over Octa are all undecidable.

About the proof

The above result on polygonal numbers does not follow directly from the Main Theorem even if we use the lemma.

Observe that

$$8T_z+1 = (2z+1)^2, \quad 3p_8(z)+1 = (3z-1)^2, \quad 24p_5(z)+1 = (6z-1)^2,$$

and hence

$$\begin{aligned}\{8t+1 : t \in \text{Tri}\} &= \{z^2 : z \in \mathbb{Z} \wedge 2 \nmid z\}, \\ \{8q+1 : t \in \text{Octa}\} &= \{z^2 : z \in \mathbb{Z} \wedge 3 \nmid z\}, \\ \{24r+1 : t \in \text{Pen}\} &= \{z^2 : z \in \mathbb{Z} \wedge 2 \nmid z \wedge 3 \nmid z\}.\end{aligned}$$

To prove the theorem, we need to modify our proof of the Main Theorem to let w (a power of p) has the form ps^2 with $2 \nmid s$, or $3 \nmid s$, or $\gcd(s, 6) = 1$. For example, to deal with triangular numbers, we take $p = 3$ and write $w = p(8t+1)$ with $t \in \text{Tri}$.

On the set of primes

Let \mathcal{P} be the set of all (positive) primes.

Matiyasevich (1975). There is a polynomial $P(x_1, \dots, x_{10}) \in \mathbb{Z}[x_1, \dots, x_{10}]$ such that

$$\mathcal{P} = \mathbb{N} \cap \{P(x_1, \dots, x_{10}) : x_1, \dots, x_{10} \in \mathbb{N}\}.$$

Theorem (Sun, arXiv:1704.03504). There are polynomials $\hat{P}(z_1, \dots, z_{20}), \tilde{P}(z_1, \dots, z_{21})$ with integer coefficients such that

$$\begin{aligned} \mathcal{P} &= \mathbb{N} \cap \{\hat{P}(z_1^2, \dots, z_{20}^2) : z_1, \dots, z_{20} \in \mathbb{N}\} \\ &= \mathbb{N} \cap \{\tilde{P}(z_1(3z_1 + 2), \dots, z_{21}(3z_{21} + 2)) : z_1, \dots, z_{21} \in \mathbb{N}\}. \end{aligned}$$

In the proof we need the Putnam trick (1969): For any polynomial $P(x) \in \mathbb{Z}[x]$, we have

$$\mathbb{N} \cap \{(x+1)(1 - P(x)^2) - 1 : x \in \mathbb{N}\} = \{x \in \mathbb{N} : P(x) = 0\}.$$

We also use the observation that any prime has the form $x^2 + y^2 + 2z^2$ (or $p_8(x) + p_8(y) + 2p_8(z)$) with $x, y, z \in \mathbb{Z}$.

References

For main sources of my work mentioned here, you may look at:

1. Z.-W. Sun, *Reduction of unknowns in Diophantine representations*, Sci. China Math. 35(1992), 257–269.
2. Z.-W. Sun, *A new relation-combining theorem and its application*, Z. Math. Logik Grundlag. Math. 38(1992), 209-212.
3. Z.-W. Sun, *Further results on Hilbert's tenth problem*, arXiv:1704.03504, <http://arxiv.org/abs/1704.03504>.

Thank you!